

How short might be the longest run in a dynamical coin tossing sequence

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Dedicated to the 100th anniversary of the birthday of Béla Gyires

Abstract. Let X_1, X_2, \dots denote i.i.d. random bits each taking the values 1 and 0 with respective probabilities $1/2$ and $1/2$. A well-known theorem of ERDŐS and RÉNYI [2] describes the limit distribution of the length of the longest contiguous run of ones in X_1, X_2, \dots, X_n as $n \rightarrow \infty$. BENJAMINI et al. ([1] Theorem 1.4) demonstrated the existence of unusual times, provided that the bits undergo equilibrium dynamics in time. In fact they prove that the dynamics produces much longer runs than the original model. In the present paper we study the length of the shortest run in the presence of the dynamics.

1. Introduction

Let X_1, X_2, \dots be a sequence of i.i.d.r.v.'s with distribution

$$\mathbf{P}\{X_1 = 1\} = \mathbf{P}\{X_1 = 0\} = \frac{1}{2}.$$

Further let

$$S_n = X_1 + X_2 + \dots + X_n, \quad I(n, a) = \max_{0 \leq k \leq n-a} (S_{k+a} - S_k), \quad (a \leq n)$$

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and Z_n be the largest integer for which $I(n, Z_n) = Z_n$, i.e. Z_n is the length of the longest run of heads in a Bernoulli trial up to n .

A characterization of the sequence $\{Z_n\}$ was given by ERDŐS and RÉVÉSZ at 1976 ([3]).

Theorem 1 (RÉVÉSZ [6], Theorem 7.2). *Assume that $\sum_{n=1}^{\infty} 2^{-a_n} < \infty$. Then*

$$Z_n \leq a_n \quad \text{a.s.} \quad (1.1)$$

if n is large enough.

Assume that $\sum_{n=1}^{\infty} 2^{-a_n} = \infty$. Then

$$Z_n \geq a_n \quad \text{a.s. i.o.} \quad (1.2)$$

$$Z_n \leq \log n - \log_3 n + \log_2 e - 1 + \varepsilon \quad \text{a.s. i.o.} \quad (1.3)$$

$$Z_n > \log n - \log_3 n + \log_2 e - 2 - \varepsilon \quad \text{a.s.} \quad (1.4)$$

if n is large enough where \log is the logarithm with base 2, \log_p is the p -th iterated logarithm, ε is an arbitrary positive number.

Slightly stronger versions of (1.3) and (1.4) are given in [4] and [7]. Note that (1.1) and (1.2) imply

$$Z_n \leq \log n + \log_2 n + (1 + \varepsilon) \log_3 n \quad \text{a.s.}$$

if n is large enough and

$$Z_n \geq \log n + \log_2 n + \log_3 n \quad \text{a.s. i.o.}$$

In [3] we also investigated the following question: what is length of the longest run containing at most one (or at most T ($T = 1, 2, \dots$)) 0's. Let $Z_n^{(T)}$ be the largest integer for which

$$I(n, Z_n^{(T)}) \geq Z_n^{(T)} - T.$$

A generalization of Theorem 1 is the following:

Theorem 2 (RÉVÉSZ [6], Theorem 7.4). *Let $\{a_n\}$ be a sequence of positive numbers and let*

$$A_T(\{a_n\}) = \sum_{n=1}^{\infty} a_n^T 2^{-a_n}.$$

Assume that

$$A_T(\{a_n\}) < \infty.$$

Then

$$Z_n^{(T)} \leq a_n \quad \text{a.s.}$$

if n is large enough.

Assume that

$$A_T(\{a_n\}) = \infty.$$

Then

$$Z_n^{(T)} \geq a_n \quad \text{a.s. i.o.}$$

Further

$$Z_n^{(T)} \leq \log n + T \log_2 n - \log_3 n - \log T! + \log_2 e - 1 + \varepsilon \quad \text{a.s. i.o.}$$

$$Z_n^{(T)} \geq \log n + T \log_2 n - \log_3 n - \log T! + \log_2 e - 2 - \varepsilon \quad \text{a.s.}$$

Given an array $\{X_n^{(j)} : j, n = 0, 1, 2, \dots\}$ of i.i.d.r.v.'s with distribution

$$\mathbf{P}\{X_0^{(0)} = 1\} = \mathbf{P}\{X_0^{(0)} = 0\} = \frac{1}{2}$$

and for each n an independent Poisson process $\{\psi_n^{(j)}\}$ ($j \geq 0$) of rate 1. Define

$$X_n(t) = X_n^{(j)} \quad \text{for} \quad \psi_n^{(j-1)} \leq t < \psi_n^{(j)} \quad (j = 1, 2, \dots)$$

where $\psi_n^{(0)} = 0$ for each n .

Let

$$S_n(t) = \sum_{k=1}^n X_k(t), \quad I(n, a, t) = \max_{0 \leq k \leq n-a} (S_{k+a}(t) - S_k(t)), \quad (0 \leq a \leq n)$$

$Z_n(t)$ be the largest integer for which $I(n, Z_n(t), t) = Z_n(t)$ i.e. $Z_n(t)$ is the length of the longest head-run at t up to n .

Similarly let $Z_n(t, T)$ be the largest integer for which $I(n, Z_n(t, T), t) \geq Z_n(t, T) - T$.

A characterization of $\{Z_n(t), n = 1, 2, \dots, 0 \leq t \leq 1\}$ is given by BENJAMINI *et al.* (2003 [1]). The properties of $\{Z_n(t, T), n = 1, 2, \dots, 0 \leq t \leq 1, T = 1, 2, \dots\}$ were studied by KHOSHNEVISAN *et al.* (2007 [5]).

Theorem 3 (BENJAMINI *et al.* [1] Theorem 1.4). Assume that $\sum_{n=1}^{\infty} a_n 2^{-a_n} < \infty$. Then

$$\forall t \in (0, 1) \text{ we have } Z_n(t) \leq a_n \quad \text{a.s. if } n \text{ is large enough.} \quad (1.5)$$

Assume that $\sum_{n=1}^{\infty} a_n 2^{-a_n} = \infty$. Then

$$\exists t \in (0, 1) \text{ such that } Z_n(t) \geq a_n \quad \text{a.s. i.o.} \quad (1.6)$$

Note that (1.5) and (1.6) imply

$$Z_n(t) \leq \log n + 2 \log_2 n + (1 + \varepsilon) \log_3 n \quad \text{a.s.}$$

for any $t \in (0, 1)$ if n is large enough and $\exists t \in (0, 1)$ such that

$$Z_n(t) \geq \log n + 2 \log_2 n + \log_3 n \quad \text{a.s. i.o.}$$

Theorem 4 (KHOSHNEVISAN *et al.* [5]). *Assume that*

$$A_{T+1}(\{a_n\}) < \infty.$$

Then

$$Z_n(t, T) \leq a_n \quad \text{a.s.}$$

for any $0 \leq t \leq 1$ if n large enough.

Assume that

$$A_{T+1}(\{a_n\}) = \infty.$$

Then there exists a t ($0 \leq t \leq 1$) such that

$$Z_n(t, T) \geq a_n \quad \text{a.s. i.o.}$$

The main goal of this paper is to prove the following:

Theorem 5. *For any $0 \leq t \leq 1 \leq 1/8$ we have*

$$Z_n(t) \geq \log n - \log_3 n - C \quad \text{a.s.}$$

if n is large enough and

$$C > 6 + \log \frac{4}{e-1}.$$

2. Combinatorial lemmas

Let $\nu(2), \nu(3), \dots, \nu(K)$ be i.i.d.r.v.'s with

$$\mathbf{P}\{\nu(2) = i\} = \frac{1}{K} \quad (i = 1, 2, \dots, K).$$

Let $X(1, 1), X(1, 2), \dots, X(1, K), Y_2, Y_3, \dots, Y_K$ be i.i.d.r.v.'s with

$$\mathbf{P}\{X(1, 1) = 1\} = \mathbf{P}\{X(1, 1) = 0\} = \frac{1}{2}$$

being independent on $\nu(2), \nu(3), \dots, \nu(K)$.

Let

$$X(2, i) = \begin{cases} X(1, i) & \text{if } \nu(2) \neq i, \\ Y_2 & \text{if } \nu(2) = i, \end{cases} \quad X(3, i) = \begin{cases} X(2, i) & \text{if } \nu(3) \neq i, \\ Y_3 & \text{if } \nu(3) = i, \end{cases}$$

$$X(k, i) = \begin{cases} X(k-1, i) & \text{if } \nu(k) \neq i, \\ Y_k & \text{if } \nu(k) = i, \end{cases} \quad \text{where } (i = 1, 2, \dots, K, k = 4, 5, \dots, K).$$

$$B(k) = \bigcap_{i=1}^K \{X(k, i) = 1\}, \quad p(K) = \mathbf{P}\left\{ \bigcup_{k=1}^K B(k) \right\}.$$

Lemma 1.

$$\frac{(e-1)K}{e^{2K+1}} \leq p(K) \leq \frac{K}{2^K}. \quad (2.1)$$

PROOF. The upper part of (2.1) is trivial. Now we turn to its lower part. Clearly

$$p(K) = \mathbf{P}\{B(1)\} + \mathbf{P}\{\overline{B(1)}B(2)\} + \mathbf{P}\{\overline{B(1)}\overline{B(2)}B(3)\} \\ + \dots + \mathbf{P}\{\overline{B(1)}\overline{B(2)}\dots\overline{B(K-1)}B(K)\},$$

(where \overline{A} is the complement of A),

$$\mathbf{P}\{B(1)\} = \frac{1}{2^K}, \quad \mathbf{P}\{\overline{B(1)}B(2)\} = \frac{1}{2^{K+1}},$$

$$\mathbf{P}\{\overline{B(1)}\overline{B(2)}B(3)\} = \mathbf{P}\{\nu(3) \neq \nu(2), X(3, \nu(3)) = X(2, \nu(2)) = 1, \\ X(1, \nu(3)) = 0, X(1, k) = 1 \text{ if } \{k \neq \nu(2) \text{ and } k \neq \nu(3)\}\} + \mathbf{P}\{\nu(3) = \nu(2), \\ X(3, \nu(3)) = 1, X(2, \nu(2)) = X(1, \nu(2)) = 0, X(1, k) = 1 \text{ if } k \neq \nu(2)\} \\ = \mathbf{P}\{\nu(3) \neq \nu(2)\} \frac{1}{2^{K+1}} + \mathbf{P}\{\nu(3) = \nu(2)\} \frac{1}{2^{K+2}}.$$

Since

$$\mathbf{P}\{\nu(2) \neq \nu(3)\} = 1 - \mathbf{P}\{\nu(2) = \nu(3)\} = 1 - \frac{1}{K}$$

we have

$$\mathbf{P}\{\overline{B(1)}\overline{B(2)}B(3)\} = \frac{1}{2^{K+1}} \left(1 - \frac{1}{2K}\right).$$

Observe that

$$\{X(1, \nu(k)) = 0, X(k, m) = 1 \text{ } (m = 1, 2, \dots, K), \nu(j) \neq \nu(k) \text{ } (j = 2, 3, \dots, k-1)\} \\ \subset \overline{B(1)}\overline{B(2)}\dots\overline{B(k-1)}B(k).$$

Since

$$\mathbf{P}\{X(1, \nu(k)) = 0, X(k, m) = 1 \ (m = 1, 2, \dots, K), \nu(j) \neq \nu(k) \\ (j = 2, 3, \dots, k-1)\} = \frac{1}{2^{K+1}} \left(1 - \frac{1}{K}\right)^{k-2},$$

we have

$$\mathbf{P}\{\overline{B(1)B(2)} \dots \overline{B(k-1)B(k)}\} \geq \frac{1}{2^{K+1}} \left(1 - \frac{1}{K}\right)^{k-2} \quad (k = 2, 3, \dots, K)$$

and

$$p(K) \geq \frac{1}{2^K} + \sum_{k=2}^K \frac{1}{2^{K+1}} \left(1 - \frac{1}{K}\right)^{k-2} = \frac{1}{2^K} + \frac{K}{2^{K+1}} \left(1 - \left(1 - \frac{1}{K}\right)^{K-1}\right) \\ \geq \frac{e-1}{e} \frac{K}{2^{K+1}}.$$

Hence Lemma 1 is proved.

Let $\mu(2), \mu(3), \dots, \mu(2K)$ be i.i.d.r.v.'s with

$$\mathbf{P}\{\mu(2) = i\} = \frac{1}{2K} \quad (i = 1, 2, \dots, 2K).$$

Let $Z(1, 1), Z(1, 2), \dots, Z(1, 2K), U_2, U_3, \dots, U_{2K}$ be i.i.d.r.v.'s with

$$\mathbf{P}\{Z(1, 1) = 1\} = \mathbf{P}\{Z(1, 1) = 0\} = \frac{1}{2}$$

being independent on $\mu(2), \mu(3), \dots, \mu(2K)$.

Let

$$Z(2, i) = \begin{cases} Z(1, i) & \text{if } \mu(2) \neq i, \\ U_2 & \text{if } \mu(2) = i, \end{cases} \quad Z(k, i) = \begin{cases} Z(k-1, i) & \text{if } \mu(k) \neq i, \\ U_k & \text{if } \mu(k) = i \end{cases}$$

$(k = 3, 4, \dots, 2K, i = 1, 2, \dots, 2K).$ □

Lemma 2. *Let*

$$q(2K) = \mathbf{P}\{\exists k, j : 1 \leq k \leq 2K, 1 \leq j \leq K \text{ such that } Z(k, j) = \dots \\ = Z(k, j + K - 1) = 1\}.$$

Then

$$\Delta \frac{K^2}{2^K} \leq q(2K) \leq \frac{K^2}{2^K} \tag{2.2}$$

where

$$\Delta = \frac{e-1}{4e^2} = 0.0581 \dots$$

PROOF. The upper part of (2.2) is trivial. Now we turn to its lower part. Let

$$A_j = \{\exists k : 1 \leq k \leq 2K, Z(k, j) = Z(k, j + 1) = \dots = Z(k, j + K - 1) = 1\},$$

$$C_j = \{Z(1, j) = Z(2, j) = \dots = Z(2K, j) = 0\} \quad (j = 1, 2, \dots, K).$$

Then

$$q(2K) = \mathbf{P}\left\{\bigcup_{j=1}^K A_j\right\} = \mathbf{P}\{A_1\} + \mathbf{P}\{\overline{A_1}A_2\} + \dots + \mathbf{P}\{\overline{A_1}\overline{A_2}\dots\overline{A_{K-1}}A_K\}$$

$$\geq \mathbf{P}\{A_1\} + \mathbf{P}\{C_1\overline{A_1}A_2\} + \dots + \mathbf{P}\{C_{K-1}\overline{A_1}\overline{A_2}\dots\overline{A_{K-1}}A_K\}$$

$$= \mathbf{P}\{A_1\} + \mathbf{P}\{C_1A_2\} + \dots + \mathbf{P}\{C_{K-1}A_K\} = \mathbf{P}\{A_1\} + (K - 1)\mathbf{P}\{C_1A_2\}.$$

Clearly

$$\{Z(1, 1) = 0, \mu_2 > 1, \mu_3 > 1, \dots, \mu_{2K} > 1\} \subset C_1,$$

$$\mathbf{P}\{C_1\} \geq \frac{1}{2} \left(1 - \frac{1}{2K}\right)^{2K-1} \geq \frac{1}{2e}.$$

Let

$$V_{2K} = \#\{k : k \leq 2K, \mu_k \leq K\}.$$

Then

$$V_{2K} \geq (1 - o(1))K \quad \text{a.s.}$$

Hence by Lemma 1 we have

$$\mathbf{P}\{A_1\} = \mathbf{P}\{A_2\} \geq \frac{e - 1}{e} \frac{K}{2^{K+1}}$$

which implies (2.2). □

3. Proof of Theorem 5

Let

$$\Psi(K) = \{\psi_n^{(j)}, n = 1, 2, \dots, 2K, j = 1, 2, \dots\}.$$

Let $\tau_1 < \tau_2 < \dots$ be the ordered elements of $\Psi(K)$ i.e. τ_1 is the smallest element of $\Psi(K)$, τ_2 is the second smallest and so on. Note $\tau = \{\tau_1, \tau_2, \dots\}$ is a Poisson process of parameter $2K$.

Define n_1 by the equation

$$\psi_{n_1}^{(1)} = \tau_1.$$

Clearly

$$\mathbf{P}\{n_1 = i\} = \frac{1}{2K} \quad (i = 1, 2, \dots, 2K).$$

Similarly define $(n_2, j_2), (n_3, j_3), \dots$ by the equations

$$\psi_{n_2}^{(j_2)} = \tau_2, \quad \psi_{n_3}^{(j_3)} = \tau_3, \dots \tag{3.1}$$

Clearly $(n_2, j_2), (n_3, j_3), \dots$ are uniquely defined a.s. by (3.1) and we have

$$\mathbf{P}\{n_\ell = i\} = \frac{1}{2K} \quad (i = 1, 2, \dots, 2K, \ell = 1, 2, \dots, 2K).$$

Introduce the following definitions:

- (i) Let $J(K)$ be the set of those j 's ($j \leq 2K$) for which there exists an $n \in [1, K]$ such that for any $t \in [\tau_j, \tau_{j+1})$ we have $X_m(t) = 1$ for each $m \in [n, n+K-1]$.
- (ii) Let $F(K)$ be the event that the set $J(K) \neq \emptyset$.
- (iii) For any $t > 0$ let $j(t)$ be the largest integer for which $\tau_j \leq t$.

If $J(K)$ consists of exactly one element then for any $j \leq K$ we have

$$\mathbf{P}\{j \in J(K)\} = \frac{1}{K}, \tag{3.2}$$

if it consists of more than one element then for any $j \leq K$ we have

$$\mathbf{P}\{j \in J(K)\} \geq \frac{1}{K}. \tag{3.3}$$

Lemma 3. *Let $j \leq K$ and $\tau_j \leq t < \tau_{j+1}$. Then we have*

$$\mathbf{P}\left\{\tau_{j(t)+1} - \tau_{j(t)} \geq \tau_{j(t)+1} - t \geq \frac{1}{2K}\right\} = \frac{1}{e}, \tag{3.4}$$

$$\mathbf{P}\{Z_{2K}(t) \geq K \mid F(K)\} = \mathbf{P}\{j(t) \in J(K) \mid F(K)\} \geq \frac{1}{K}. \tag{3.5}$$

Consequently

$$\mathbf{P}\left\{Z_{2K}(t) \geq K, \tau_{j(t)+1} - t \geq \frac{1}{2K} \mid F(K)\right\} \geq \frac{1}{eK}. \tag{3.6}$$

PROOF. Since τ is a Poisson process of parameter $1/2K$, we have (3.4). (3.5) follows from (3.2) and (3.3). (3.4) and (3.5) combined imply (3.6). \square

Lemma 4.

$$\mathbf{P}\{F(K)\} \geq \Delta K^2 2^{-K}. \tag{3.7}$$

PROOF. It is a trivial consequence of Lemma 2. □

Lemma 5. *Let $j \leq K$ and $\tau_j \leq t < \tau_{j+1}$. Then we have*

$$\mathbf{P}\left\{\tau_{j(t)+1} - t \geq \frac{1}{2K}, Z_{2K}(t) \geq K\right\} \geq \frac{\Delta}{e} K 2^{-K}. \tag{3.8}$$

PROOF. Clearly the r.v.'s $Z_{2K}(t)$ and $\tau_{j(t)+1} - \tau_{j(t)}$ are independent. Hence

$$\begin{aligned} & \mathbf{P}\left\{\tau_{j(t)+1} - t \geq \frac{1}{2K}, Z_{2K}(t) \geq K\right\} \\ &= \mathbf{P}\left\{\tau_{j(t)+1} - t \geq \frac{1}{2K}, Z_{2K}(t) \geq K, F(K)\right\} \\ &= \mathbf{P}\left\{\tau_{j(t)+1} - t \geq \frac{1}{2K}\right\} \mathbf{P}\{Z_{2K}(t) \geq K \mid F(K)\} \mathbf{P}\{F(K)\}. \end{aligned}$$

Hence we have (3.8) by (3.4), (3.5) and (3.6). □

Introduce two further definitions:

(iv) Let $Z(\ell, K, t)$ be the largest integer for which there exists an integer

$$b \in [2\ell K, (2\ell + 1)K)$$

such that

$$S_{b+K}(t) - S_b(t) = Z(\ell, K, t)$$

that is $Z(\ell, K, t)$ is the length of the longest head-run at

$$t \in [2\ell K, (2\ell + 2)K).$$

(v) Define $\Psi(\ell, K)$, $\tau_i(\ell)$, $F(\ell, K)$, $j(\ell, t)$ just like above using the block

$$[2\ell K, (2\ell + 2)K)$$

instead of the block $(0, 2K)$.

Let

$$Q := Q(K, N) = \#\{\ell : \ell \leq [N/2K] \text{ for which } F(\ell, K) \text{ holds true}\}$$

and let

$$K = \log N - \log_3 N - C.$$

Lemma 6.

$$\mathbf{P}\{Q(K, N) \leq \Delta 2^{C-2} \log N \log_2 N\} \leq N^{-O(1) \log_2 N}. \quad (3.9)$$

PROOF. Apply the large deviation theorem (RÉVÉSZ [6], Theorem 2.3) with

$$n = \lfloor N/2K \rfloor, \quad p = \frac{\Delta K^2}{2K}, \quad \varepsilon = \frac{p}{2}.$$

Since

$$\frac{np}{2} = \Delta 2^{C-2} \log N \log_2 N, \quad \frac{n\varepsilon^2}{2pq(1 + \frac{\varepsilon}{2pq})^2} = O(1) \log N \log_2 N,$$

we have (3.9). \square

Let $D(t) = D(t, N, K)$ be the event that there exists an $\ell = \ell(t)$ ($0 \leq \ell \leq \lfloor N/2K \rfloor$) for which $Z(\ell, K, t) \geq K$ and $\tau_{j(t)+1}(\ell) - t \geq 1/2K$.

Lemma 7.

$$\mathbf{P}\{\overline{D(t)}\} \leq \frac{2}{(\log N)^{\Delta 2^{C-2} e^{-1}}}. \quad (3.10)$$

PROOF. By (3.6) and Lemma 6 we have

$$\begin{aligned} bp\{\overline{D(t)}\} &= \mathbf{E}\mathbf{P}\{\overline{D(t)} \mid Q\} \geq \mathbf{E}\left(1 - \frac{1}{2K}\right)^Q \\ &= \mathbf{E}\left(\left(1 - \frac{1}{eK}\right)^Q \mid Q \leq \Delta 2^{C-2} \log N \log_2 N\right) \mathbf{P}\{Q \leq \Delta 2^{C-2} \log N \log_2 N\} \\ &\quad + \mathbf{E}\left(\left(1 - \frac{1}{eK}\right)^Q \mid Q > \Delta 2^{C-2} \log N \log_2 N\right) \mathbf{P}\{Q > \Delta 2^{C-2} \log N \log_2 N\} \\ &\leq \mathbf{P}\{Q \leq \Delta 2^{C-2} \log N \log_2 N\} + \left(1 - \frac{1}{eK}\right)^{\Delta 2^{C-2} \log N \log_2 N} \leq \\ &\leq N^{-O(1) \log_2 N} + \exp\left(-\frac{\Delta 2^{C-2}}{e} \log_2 N\right) \leq 2(\log N)^{-\Delta 2^{C-2} e^{-1}}. \end{aligned}$$

Hence we have (3.10). \square

Lemma 8.

$$\mathbf{P}\left\{\tau_{2K} < \frac{1}{8}\right\} \leq e^{-BK}. \quad (3.11)$$

Consequently

$$\mathbf{P}\left\{\min_{\ell \leq \lfloor N/2K \rfloor} \tau_{2K}(\ell) < \frac{1}{8}\right\} \leq e^{-BK} \quad (3.12)$$

for a $B > 0$.

PROOF. Since

$$\frac{d}{dx} \mathbf{P}\{2K\tau_K < x\} = \frac{x^{K-1}e^{-x}}{(K-1)!},$$

by an easy calculation we have (3.11). □

Lemma 9. *The event*

$$\bigcap_{m=1}^{2K} D\left(\frac{m}{2K}\right)$$

implies that for any $m \leq 2K$ there exists an $\ell \leq [N/2K]$ such that

$$\tau_{j(m/2K)+1}(\ell) - \frac{m}{2K} \geq \frac{1}{2K}$$

and

$$Z(\ell, K, t) \geq K$$

if

$$t \leq \min_{\ell \leq [N/2K]} \tau_{2K}(\ell).$$

PROOF. It follows straight from the definition of $D(\cdot)$. □

PROOF OF THEOREM 5. By Lemma 7 we have

$$\begin{aligned} \mathbf{P}\left\{\bigcap_{j=1}^{2K} D\left(\frac{j}{2K}\right)\right\} &= 1 - \mathbf{P}\left\{\bigcup_{j=1}^{2K} \overline{D\left(\frac{j}{2K}\right)}\right\} \geq 1 - 2K\mathbf{P}\{\overline{D(t)}\} \\ &\geq 1 - 2K2(\log N)^{\Delta 2^{C-2}e^{-1}} \geq 1 - 4(\log N)^{-\Delta 2^{C-2}e^{-1}+1}. \end{aligned}$$

Let

$$C > 3 + \log \frac{e}{\Delta}$$

and

$$N_m = 2^m \quad (m = 1, 2, \dots).$$

Then

$$\Delta 2^{C-2}e^{-1} - 1 > 1$$

and the event

$$\bigcap_{j=1}^{2K} D\left(\frac{j}{2K}\right) \quad (K = K_m = \log N_m - \log_3 N_m - C)$$

occurs with probability 1 if m is large enough. Consequently for any $t \leq 1/8$

$$Z_{N_m}(t) \geq \log N_m - \log_3 N_m - C$$

if m is large enough which easily implies Theorem 5. □

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