

## An algorithm determining cycles of polynomial mappings in integral domains

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**Abstract.** In the first part of this paper we show how all normalized cycles could be found in a domain  $R$ , provided all nontrivial solutions in units of  $u + v = 1$  and  $u + v + w = 1$  are given. Then we give an effective method to find all normalized cycles in the ring of integers  $Z_K$  in any algebraic number field  $K$ . Finally, we deal with polynomial orbits.

For a commutative ring  $S$  with unity a tuple  $(x_0, x_1, \dots, x_{n-1})$  of distinct elements from  $S$  is called a (*polynomial*) *cycle* if for some  $f \in S[X]$  we have  $f(x_0) = x_1, f(x_1) = x_2, \dots, f(x_{n-2}) = x_{n-1}, f(x_{n-1}) = x_0$ . The number  $n$  is called the *length* of this cycle. A cycle  $x_0, x_1, \dots$  is called *normalized* provided  $x_0 = 0, x_1 = 1$ .

In the first section we show how all normalized cycles could be found in a domain  $R$ , provided all nontrivial solutions in units of  $u + v = 1$  and  $u + v + w = 1$  are given. In the second section we give an effective method to find all normalized cycles in the ring of integers  $Z_K$  in any algebraic number field  $K$ . In the last section we deal with polynomial orbits.

### 1. A usefulness of $u + v = 1$ and $u + v + w = 1$ in units

In [H-KNa2] the following theorem was established:

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**Theorem 0.** *Let  $R$  be an integral domain and assume that for every nonzero  $b \in R$  each of the equations:  $x_1 + bx_2 = 1$ ,  $b(x_1 + x_2) + x_3 = 1$ ,  $x_1 + x_2 + x_3 + x_4 + x_5 = 1$  has only finitely many nontrivial solutions  $x_i \in R^\times$ . Then there are only finitely many normalized cycles in  $R$  of a given length.*

Let  $R$  be a commutative domain in which the equations  $u + v = 1$  and  $u + v + w = 1$  have only finitely many solutions in units  $\neq 1$  (this assumption is satisfied for finitely generated domains of 0 characteristics). Let us define  $\mathcal{A}$  as the set of all solutions  $(\alpha_i, \beta_i, \gamma_i)$ ,  $i = 1, \dots, m$  of  $\alpha + \beta + \gamma = 1$  with invertible  $\alpha, \beta, \gamma$  distinct from 1, and  $|\mathcal{A}| = m$ . We also define  $\mathcal{B}$  as the set of all solutions  $(\delta_j, \epsilon_j)$ ,  $j = 1, \dots, m_1$  of  $\delta + \epsilon = 1$  with invertible  $\delta, \epsilon$ , and  $|\mathcal{B}| = m_1$ . In the following theorem we improve Theorem 0.

**Theorem 1.** *Let  $R, \mathcal{A}, \mathcal{B}$  be as above.*

(i) *The lengths of cycles in  $R$  are bounded by  $6(m+2)^2$ .*

(ii) *Fix  $n \leq 6(m+2)^2$  and  $n \neq 4$ . We define a family of sets  $\mathcal{X}_j$  as follows.*

*Put  $\mathcal{X}_1 = \{1\}$ . For odd  $n$  we put  $\mathcal{X}_2 = \{1 - \delta_j : j = 1, \dots, m_1\}$ . For even  $n \neq 6, 12$  we put  $\mathcal{X}_2 = \{1 - \alpha_i : i = 1, \dots, m\}$ . For  $n = 12$  we put  $\mathcal{X}_2 = \{1 - \alpha_i, 1 + \alpha_i : i = 1, \dots, m\}$ . For  $n = 6$  we put  $\mathcal{X}_2 = \{1 - \alpha_i : i = 1, \dots, m\} \cup \{1 - \xi, 1 - \xi^2\}$ , where  $\xi \in R$  is a primitive third root of unity (if it exists in  $R$ , otherwise we skip the second component).*

*Having defined  $\mathcal{X}_1, \mathcal{X}_2$ , we define inductively  $\mathcal{X}_i$  for  $i \geq 3$  by*

*$\mathcal{X}_i = \{a(x - y) + y : a \in \mathcal{X}_2, x \in \mathcal{X}_{i-1}, y \in \mathcal{X}_{i-2}, x - y \text{ is invertible}\}$ .*

*Then for  $n \neq 4, 6$  any normalized cycle  $(x_0 = 0, x_1 = 1, x_2, \dots, x_{n-1})$  of length  $n$  satisfies  $x_i \in \mathcal{X}_i$ .*

*Any normalized cycle  $(x_0 = 0, x_1 = 1, x_2, \dots, x_5)$  of length 6 satisfies  $x_i \in \mathcal{X}_i$ , except for  $\text{char } R = 3$ , where, in addition,  $(0, 1, 1 - u, 2 - u, u - 1, u)$  is a cycle for any invertible  $u \neq 1, 2$ .*

(iii) *Any normalized cycle in  $R$  of length 4 is of one of the following forms:*

(a)  $(0, 1, 1 - \alpha_i, \beta_i)$ , where  $1 \leq i \leq m$  and the ratio  $(1 - \alpha_i)/(1 - \beta_i)$  is invertible;

(b)  $(0, 1, 1 + \alpha_i, \alpha_i)$ , where  $1 \leq i \leq m$  and the ratio  $(1 + \alpha_i)/(1 - \alpha_i)$  is invertible;

(c)  $(0, 1, 1 + \epsilon, \epsilon)$ , where  $\epsilon$  satisfies  $\epsilon^2 + 1 = 0$ ;

(d) (only for  $\text{char } R = 2$ )  $(0, 1, 1 + v, v)$ , where  $v$  is any unit  $\neq 1$ .

**PROOF.** (i) Let  $(0, 1, x_2, \dots, x_{n-1})$  be a cycle. Lemma 1 from [Na1] gives that for any  $1 \leq k \leq n-1$  satisfying  $(k(k-2), n) = 1$  the elements  $x_k, x_1 - x_2 = 1 - x_2$

and  $x_2 - x_k$  are invertible. Thus for such  $k$  the triples  $(\alpha, \beta, \gamma) = (1 - x_2, x_2 - x_k, x_k)$  are distinct solutions of the 3-unit equation  $\alpha + \beta + \gamma = 1$ . Among them there are at most two trivial solutions, i.e. when  $k = 1$  or  $x_2 - x_k = 1$ . Thus the number  $l$  of integers  $k \in [1, n - 1]$  satisfying  $(k(k - 2), n) = 1$  cannot exceed  $m + 2$ . Since

$$l = \begin{cases} n \prod_{p|n} \left(1 - \frac{2}{p}\right) & \text{if } 2 \nmid n, \\ \frac{n}{2} \prod_{2 \neq p|n} \left(1 - \frac{2}{p}\right) & \text{if } 2 \mid n, \end{cases} \tag{1}$$

and  $p^\alpha(1 - 2/p) \geq \sqrt{p^\alpha}$  for prime  $p \geq 5$  and  $\alpha \geq 1$ , we get  $m + 2 \geq l \geq \sqrt{n}/\sqrt{6}$ . Thus  $n \leq 6(m + 2)^2$ .

(ii) Let  $n \geq 3$  and let  $(0, 1, x_2, \dots, x_{n-1})$  be a cycle in  $R$ . Our first aim is to prove that  $x_2 \in \mathcal{X}_2$ .

If  $n$  is odd, then  $1 - x_2 = x_1 - x_2$  and  $x_2$  are invertible. Thus  $x_2 = 1 - \delta_j$  for some  $1 \leq j \leq m_1$ . Hence  $x_2 \in \mathcal{X}_2$ .

Assume now that  $n \neq 4, 6, 12$  is even. Then the number  $l$  from (1) satisfies  $l \geq 3$ . So there exists  $k \in [2, n - 1]$  such that  $(k(k - 2), n) = 1$  and  $x_2 - x_k \neq 1$ . Hence  $(1 - x_2, x_2 - x_k, x_k) = (\alpha_i, \beta_i, \gamma_i)$  for some  $1 \leq i \leq m$ , and  $x_2 \in \mathcal{X}_2$  follows.

For  $n = 12$  the triples  $(1 - x_2, x_2 - x_7, x_7)$  and  $(1 - x_6, x_6 - x_7, x_7)$  are solutions of the 3-unit equation  $\alpha + \beta + \gamma = 1$ . If the first solution is not trivial, then  $1 - x_2 = \alpha_i$ ,  $x_2 = 1 - \alpha_i$  for some  $1 \leq i \leq m$ . Otherwise  $x_2 - x_7 = 1$ , but then the second solution is not trivial, and  $x_7 = \alpha_i$ ,  $x_2 = 1 + \alpha_i$  for some  $i$ . Hence  $x_2 \in \mathcal{X}_2$ .

Assume that  $n = 6$ . Put  $x_6 = 0$ ,  $x_7 = 1$  and  $y_i = x_i - x_{i-1}$  for  $i \in [2, 7]$ . Hence  $y_2, \dots, y_7$  are invertible. Put  $y_2 = -u$ . Assume that  $x_2 \neq 1 - \alpha_i$  (i.e.  $u \neq \alpha_i$ ) for  $1 \leq i \leq m$ .

Take any  $i \in [3, 7]$ . Then  $x_2 \sim x_i - x_{i-2} = y_{i-1} + y_i$ , and  $1 - u = x_2 = \delta(y_{i-1} + y_i)$  for some invertible  $\delta$ . Thus  $(u, \delta y_{i-1}, \delta y_i)$  is the trivial solution of the equation  $\alpha + \beta + \gamma = 1$  in units. Since  $u \neq 1$ , we get  $\delta y_{i-1} = 1$  or  $\delta y_i = 1$ , and  $y_i/y_{i-1} \in \{-u, -u^{-1}\}$  follows. Since  $y_7 = 1$ , we get  $y_6 \in \{-u, -u^{-1}\}$ .

If  $u = -1$ , then  $x_2 = 2$ ,  $x_3 = 3$  and  $0 = x_6 = 2 \cdot 3$ . This gives  $2 = 0$  or  $3 = 0$ , i.e.  $x_2 = 0$  or  $x_3 = 0$ , a contradiction. Thus we get  $u \neq \pm 1$ .

For  $i = 2, \dots, 6$  put  $y_i = (-u)^{a_i}$ , with  $a_2 = 1$  and  $|a_i - a_{i-1}| = 1$  for  $i = 3, \dots, 6$ .

In this way we obtain 16 possibilities for the quadruple  $(a_3, a_4, a_5, a_6)$ .

There are 9 possibilities for  $(a_3, a_4, a_5, a_6)$  such that  $(a_3, a_4, a_5, a_6) \neq (0, 1, 0, 1)$  and  $a_6 \in \{\pm 1\}$ . In each of these possibilities the condition that  $0, 1, x_2, \dots, x_5$  are distinct is not satisfied. A typical such possibility is  $(a_3, a_4, a_5, a_6) = (0, 1, 0, -1)$ .

In this case  $0 = x_6 = 1 - u + 1 - u + 1 - u^{-1} = (1 - u)(1 + 1 - u^{-1})$ , and  $1 + 1 - u^{-1} = 0$  follows. This gives  $x_4 - x_1 = -u + 1 - u = -u(1 + 1 - u^{-1}) = 0$ , a contradiction.

There are 5 possibilities for  $(a_3, a_4, a_5, a_6)$  such that  $a_6 \in \{\pm 3\}$ . In any such possibility  $y_6 \in \{-u, -u^{-1}\}$  gives  $(-u)^2 = 1$  or  $(-u)^4 = 1$ . Since we have already excluded  $u = \pm 1$ , we must have  $u^2 + 1 = 0$ , and this gives a contradiction. Take for example  $(a_3, a_4, a_5, a_6) = (0, -1, -2, -3)$ . Since  $0 = x_6 = 1 - u + 1 - u^{-1} + u^{-2} - u^{-3} = (1 - u)(1 - u^{-1} - u^{-3})$ , we obtain  $0 = 1 - u^{-1} - u^{-3} = 1 - u^{-3}(1 + u^2) = 1$ , a contradiction. In other four cases we proceed in a similar manner.

Let us consider  $(a_3, a_4, a_5, a_6) = (2, 3, 4, 5)$ . Since  $0 = x_6 = 1 - u + u^2 - u^3 + u^4 - u^5 = (1 - u)(1 - u + u^2)(1 + u + u^2)$  and  $x_3 = 1 - u + u^2 \neq 0$ , we obtain  $1 + u + u^2 = 0$ . Hence  $u$  is a primitive third root of unity and  $x_2 = 1 - u \in \mathcal{X}_2$ .

Finally, let  $(a_3, a_4, a_5, a_6) = (0, 1, 0, 1)$ . Then  $0 = x_6 = 3(1 - u)$  and  $\text{char } R = 3$  follows. If an invertible  $u \neq 1, 2$  and  $\text{char } R = 3$ , then  $(0, 1, 1 - u, 2 - u, 2 - 2u, -2u)$  is a cycle for  $f(X) = 1 - uX - (u + 1)/uX(X - 1) + 1/uX(X - 1)(X - (1 - u))$ .

**Lemma 1.** *For  $n \geq 3$ ,  $n \neq 4$  let  $(0, 1, x_2, \dots, x_{n-1})$  be a cycle in  $R$  for  $f(X) \in R[X]$ . We extend the indices putting  $x_n = x_0 = 0$ ,  $x_{n+1} = x_1 = 1$ ,  $x_{n+2} = x_2$ ,  $x_{n+3} = x_3$  and so on. Assume that for some  $2 \leq r \leq n + 1$  we have  $1 - y := (x_r - x_{r-2})/(x_{r-1} - x_{r-2}) \notin \mathcal{X}_2$ . Then  $x_2 \notin \mathcal{X}_2$ .*

**PROOF.** We have that  $(0, 1, 1 - y, (x_{r+1} - x_{r-2})/(x_{r-1} - x_{r-2}), \dots, (x_{r+n-3} - x_{r-2})/(x_{r-1} - x_{r-2}))$  is a cycle for  $g(X) = (x_{r-1} - x_{r-2})^{-1}(f((x_{r-1} - x_{r-2})X + x_{r-2}) - x_{r-2}) \in R[X]$ . By what we have already proved, we have  $n = 6$ ,  $\text{char } R = 3$  and  $(x_{r+1} - x_{r-2})/(x_{r-1} - x_{r-2}) = 2 - y, \dots, (x_{r-3} - x_{r-2})/(x_{r-1} - x_{r-2}) = y$ . In particular the triple  $((x_0 - x_{r-2})/(x_{r-1} - x_{r-2}), (x_1 - x_{r-2})/(x_{r-1} - x_{r-2}), (x_2 - x_{r-2})/(x_{r-1} - x_{r-2}))$  is of one of the following forms  $(0, 1, 1 - y), (1, 1 - y, 2 - y), \dots, (y - 1, y, 0), (y, 0, 1)$ . This easily gives  $x_2 \in \{1 - y, 1 - 1/y\}$ . Since  $1 - y \in \mathcal{X}_2$  if and only if  $1 - 1/y \in \mathcal{X}_2$ , we are done.  $\square$

Summing up, for  $n \geq 3$ ,  $n \neq 4$  we showed that  $x_2 \in \mathcal{X}_2$ , except for one family of exceptions in  $\text{char } R = 3$ . Using Lemma 1, by simple induction we obtain  $x_j \in \mathcal{X}_j$  for  $j \geq 2$  provided  $x_2 \in \mathcal{X}_2$ . If  $x_2 \notin \mathcal{X}_2$ , then  $\text{char } R = 3$  and  $(0, 1, x_2, \dots, x_5)$  is of the form  $(0, 1, 1 - u, 2 - u, 2 - 2u, -2u)$ , with invertible  $u \neq 1, 2$ .

(iii) Let  $(0, 1, x_2, x_3)$  be a normalized cycle. We see that  $1 - x_2, x_3$  and  $x_3 - x_2$  are invertible and  $(1 - x_2, x_3, x_2 - x_3)$  is a solution of the 3-unit equation. If this solution is not trivial, then  $(1 - x_2, x_3, x_2 - x_3) = (\alpha_i, \beta_i, \gamma_i)$  for some  $1 \leq i \leq m$ , and  $(0, 1, x_2, x_3) = (0, 1, 1 - \alpha_i, \beta_i)$  follows.

Otherwise  $x_2 - x_3 = 1$ , and  $(0, 1, x_2, x_3) = (0, 1, 1 + v, v)$  for some unit  $v$ . Since  $x_3 - 1 \sim x_2$ , we obtain that  $1 - v = \delta(1 + v)$  for some invertible  $\delta$ . We see that  $(v, \delta, \delta v)$  is a solution of the 3-unit equation. If this solution is not trivial, then  $(0, 1, x_2, x_3) = (0, 1, 1 + \alpha_i, \alpha_i)$  for some  $1 \leq i \leq m$ . If this solution is trivial, then  $\delta = 1$ ,  $\text{char } R = 2$  or  $\delta = 1/v$ ,  $v^2 + 1 = 0$ . One easily sees that  $(0, 1, 1 + v, v)$  is a cycle if  $\text{char } R = 2$  and  $v \neq 1$  is invertible.  $\square$

*Remark 1.* Theorem 1 gives (except for two families of cycles if  $\text{char } R$  equals 2 or 3) a finite list of tuples which may be cycles. To check whether a given  $(0, 1, x_2, \dots, x_{n-1})$  is a cycle we should calculate the Lagrange interpolation polynomial realizing this cycle, and check whether its coefficients lie in  $R$ .

*Remark 2.* Generally speaking, the numbers of elements of  $\mathcal{X}_2, \mathcal{X}_3, \dots$  grow quite rapidly. In some cases one can shrink quite a lot the sets of possible values for  $x_j$ 's. For example if  $n$  is odd,  $(0, 1, x_2, \dots, x_{n-1})$  is a cycle and  $1 \leq j \leq n - 1$  satisfies  $(j(j-1), n) = 1$ , then  $(x_j, 1 - x_j) \in \mathcal{B}$ . One may also restrict the possible values for  $x_j$ 's taking into account for example that  $x_2 \sim x_3 - x_1 \sim x_4 - x_2$  and some other similar relations.

*Remark 3.* If  $(x_0, x_1, \dots, x_{n-1})$  is a cycle, then  $(0, 1, (x_2 - x_0)/(x_1 - x_0), (x_3 - x_0)/(x_1 - x_0), \dots, (x_{n-1} - x_0)/(x_1 - x_0))$  is a normalized cycle of the same length. Thus having found all normalized cycles we find the set  $\mathcal{CYCL}(R)$  of all cycle lengths in  $R$ .

*Remark 4.* From the proof of Theorem 1 we infer that for odd  $n$  all normalized cycles of length  $n$  can be found using solely the solutions of the 2-unit equation  $u + v = 1$ . Using the ideas from the proof of Theorem 1(i), we may show that the odd lengths of cycles are bounded by  $C(m_1 + 1)(\log \log(m_1 + 3))^2$  for some constant  $C$ . Thus finding all normalized cycles of even lengths is much complicated than those of odd lengths.

*Remark 5.* Theorem 1 does not lead directly to the determination of  $\mathcal{CYCL}(R)$  in the case when  $R = Z_K$  is the ring of integers of an algebraic number field  $K$ , as there is no known procedure to find all solutions of the equation  $u + v + w = 1$  in units  $\neq 1$ . An exception is formed by fields with unit rank  $\leq 1$ , where all cycle lengths were determined ([Bo] and [Ba]) for quadratic fields, [Na2] for complex cubic fields, and [Pe2] for totally complex quartic fields.

## 2. An algorithm determining all normalized cycles

Nevertheless, we have found a finitary procedure working in all number fields  $K$ , which finds all normalized cycles in  $Z_K$ , and therefore also  $\mathcal{CYCL}(Z_K)$ . It is based on some known algorithms from algebraic number theory. In the proof of Theorem 2 below we propose such a procedure.

**Theorem 2.** *There is an effective procedure, which for a given number field  $K$  finds all normalized cycles in  $Z_K$ . This procedure also finds  $\mathcal{CYCL}(Z_K)$ .*

PROOF. For any number field  $K$  the following things can be effectively calculated:

(A) the degree  $[K : Q]$ , the discriminant  $\text{disc}(K)$ , the regulator  $\text{reg}(K)$ , the class number  $h_K$ , an integral basis, a fundamental system of units, all roots of unity lying in  $K$ . One may effectively check whether a given element from  $K$  lies in  $Z_K$ .

(B) For any nonzero  $\alpha, \beta, \gamma \in Z_K$  let us define  $\mathcal{T}_K(\alpha, \beta; \gamma) = \{(u, v) : \alpha u + \beta v = \gamma, \text{ and } u, v \text{ are invertible}\}$ . Then  $\mathcal{T}_K(\alpha, \beta; \gamma)$  is finite and may be effectively found (see [S], [Gy], [EGST]).

(C) For any nonzero  $a \in Z_K$  one can effectively find the set of all (up to associates) divisors of  $a$ . Since for any  $b \mid a$  one has  $N_{K/Q}(b) \mid N_{K/Q}(a)$ , this may be completed by solving a suitable norm form equation. For the solvability of norm form equations in an effective way see [BSh], [Ga].

**Lemma 2.** *Let  $K$  be a number field. Let  $a_1, a_2, b_1, b_2, c_1, c_2 \in Z_K$  be given and satisfy  $a_1, a_2, c_1, c_2, b_1a_2 - b_2a_1 \neq 0$ . Then one can effectively determine a finite set  $\mathcal{A}_1$ , depending on  $K$  and  $a_1, a_2, \dots, c_2$ , consisting of all integers  $u \in Z_K$  satisfying  $a_i u + b_i \mid c_i$  for  $i = 1, 2$ .*

PROOF. For  $i = 1, 2$  one effectively finds a finite set  $\mathcal{D}_i$  of all (up to associates) divisors of  $c_i$ . Hence  $a_1 u + b_1 = d_1 \delta_1$ ,  $a_2 u + b_2 = d_2 \delta_2$  for some  $d_i \in \mathcal{D}_i$  and invertible  $\delta_1, \delta_2$ . This gives  $(\delta_1, \delta_2) \in \mathcal{T}_K(d_1 a_2, -d_2 a_1; b_1 a_2 - b_2 a_1)$ , so for fixed  $d_1, d_2$  we effectively find possible  $u$ . Since  $d_i$  lies in the finite set  $\mathcal{D}_i$ , we are done.  $\square$

Let  $K$  be a fixed number field. Put  $N = [K : Q]$ . Let  $(r, s)$  be the signature of  $K$ . Let  $\zeta_M = \exp(2\pi i/M)$  be the generator of the group of roots of unity lying in  $K$  and let  $\eta_1, \dots, \eta_{r+s-1}$  be any fundamental system of units in  $K$ .

Let us define  $\mathcal{S}(n)$  as the set of all normalized cycles in  $Z_K$  of length  $n$ . In order to prove the assertion it suffices to bound effectively the  $n$ 's such that  $\mathcal{S}(n)$  is non-empty, and for  $n$  less than this bound to find effectively  $\mathcal{S}(n)$ . According

to Remark 1, in order to check whether  $(0, 1, x_2, \dots, x_{n-1})$  is a cycle it suffices to check whether the (unique) polynomial  $h(X)$  of degree  $\leq n - 1$  realizing this cycle, with coefficients in  $K$ , has all its coefficients in  $Z_K$ . The polynomial  $h(X)$  is calculated by the Lagrange interpolation formula, and one may effectively check whether all its coefficients lie in  $Z_K$ .

*Remark 6.* Let  $B(R)$  be the biggest element of  $\mathcal{CYCL}(R)$ . It is known that  $B(Z_K)$  is bounded from above by some explicit expression depending on  $N = [K : \mathbb{Q}]$ . The first such estimation was given in [Na1], where  $B(Z_K)$  is bounded from above by some double exponential function in  $N$ . It was improved in [Pe1], where  $B(Z_K) \leq 2^{N+1}(2^N - 1)$  was established.

For any odd  $n \leq 2^{N+1}(2^N - 1)$  we can find effectively all elements from  $\mathcal{S}(n)$ , as explained in Remark 4, since by (B) all solutions of the 2-unit equation  $u + v = 1$  in  $Z_K$  can be effectively computed.

The procedure of finding all  $\mathcal{S}(k)$  will be completed provided for all  $n \leq 2^N(2^N - 1)$  we can effectively find  $\mathcal{S}(2n)$  having at our disposal the finite set  $\mathcal{S}(n)$ .

*Remark 7.* Assume that  $(0, 1, x_2, \dots, x_{n-1})$  is a cycle in a domain  $R$  for a polynomial  $f(X) = c_0 + c_1X + \dots + c_{n-1}X^{n-1}$ . Take any nonzero  $a \in R$ . The Lagrange interpolation polynomial for the sequence  $\xi_a = (0, a, ax_2, \dots, ax_{n-1})$  equals  $af((1/a)X)$ . Thus  $\xi_a$  is a cycle in  $R$  if and only if  $a \mid c_2, a^2 \mid c_3, \dots, a^{n-2} \mid c_{n-1}$ .

*Remark 8.* If  $(0, 1, y_2, y_3, \dots, y_{2n-1}) \in \mathcal{S}(2n)$  is a cycle for  $F(X)$ , then  $(0, 1, y_4/y_2, y_6/y_2, \dots, y_{2n-2}/y_2)$  is a cycle of length  $n$  for  $(1/y_2)(F \circ F)(y_2X) \in R[X]$ .

Owing to the last remark in order to find  $\mathcal{S}(2n)$  it suffices to find effectively for a fixed  $(0, 1, x_2, \dots, x_{n-1}) \in \mathcal{S}(n)$  all  $(0, 1, y_2, \dots, y_{2n-1}) \in \mathcal{S}(2n)$  such that  $y_{2k}/y_2 = x_k$  for all  $2 \leq k \leq n - 1$  (informally speaking  $(0, 1, x_2, \dots, x_{n-1})$  is proportional to  $(0, y_2, y_4, \dots, y_{2n-2})$ ). Let us call such  $(0, 1, y_2, \dots, y_{2n-1})$  *connected* to  $(0, 1, x_2, \dots, x_{n-1})$ .

**Lemma 3.** *Let  $n$  be given and assume that the set  $\mathcal{S}(n)$  is explicitly known. If for each sequence  $\xi = (0, 1, x_2, \dots, x_{n-1}) \in \mathcal{S}(n)$  one can effectively construct a finite set  $\mathcal{Y} = \mathcal{Y}(\xi, K)$  such that every cycle  $\eta = (0, 1, y_2, \dots, y_{2n-1})$ , connected to  $\xi$  satisfies  $y_2 \in \mathcal{Y}$ , then there exists an effective procedure to determine  $\mathcal{S}(2n)$ .*

*Such a construction exists, provided one can either effectively find a nonzero  $b \in Z_K$  such that each cycle  $\eta$  connected to  $\xi$  satisfies  $y_2 \mid b$ , or one can find*

effectively nonzero  $b, c \in Z_K$  with  $b \neq 1$  such that each cycle  $\eta$  connected to  $\xi$  satisfies  $y_2b - 1 \mid c$ .

PROOF. Suppose that we found such  $\mathcal{Y}$ . Fix any  $a \in \mathcal{Y}$ . If  $\eta$  is a cycle with  $y_2 = a$ , connected to  $\xi$ , then  $y_{2k} = ax_k$  for all  $2 \leq k \leq n - 1$ . Thus  $y_2, y_4, y_6, \dots, y_{2n-2}$  are uniquely determined by  $a$  and  $\xi$ . Let  $\mathcal{T}_K(1, 1; a) = \{(u_1, v_1), \dots, (u_t, v_t)\}$ .

Take any  $1 \leq k \leq n - 1$ , and consider  $y_{2k+1}$ . Then  $y_{2k+2} - y_{2k+1}, y_{2k+1} - y_{2k}, x_{k+1} - x_k$  are units, and  $((y_{2k+2} - y_{2k+1})/(x_{k+1} - x_k), (y_{2k+1} - y_{2k})/(x_{k+1} - x_k)) \in \mathcal{T}_K(1, 1; a)$ . This gives  $y_{2k+1} = ax_k + (x_{k+1} - x_k)v_i$  for some  $1 \leq i \leq t$ . Thus we have only finitely many and effectively computable possibilities for the values  $y_3, y_5, \dots, y_{2n-1}$ . Having a finite number of possibilities for cycles connected to  $(0, 1, x_2, \dots, x_{n-1})$  we pick those which are in fact cycles.

Let  $b \neq 0$  be effectively computable and suppose that for all cycles  $\eta$  connected to  $(0, 1, \dots, x_{n-1})$  we have  $y_2 \mid b$ . Such a set  $\mathcal{Y}$  exists by Lemma 2 in view of  $y_2 - 1 \mid 1$ .

Let  $b \neq 0, 1; c \neq 0$  be effectively computable and suppose that for all cycles  $\eta$  connected to  $\xi$  we have  $y_2b - 1 \mid c$ . Such a set  $\mathcal{Y}$  exists by Lemma 2 in view of  $y_2 - 1 \mid 1$ . □

Let  $\xi = (0, 1, x_2, \dots, x_{n-1}) \in \mathcal{S}(n)$  be fixed. Let  $\eta = (0, 1, y_2, y_3, \dots, y_{2n-1})$  be connected to  $\xi$ . Put  $y_2 = a$ . We will show two ways of finding a set  $\mathcal{Y}$  fulfilling the condition of Lemma 3. The sets  $\mathcal{Y}$  obtained in the two ways below may differ.

**First way.** Let  $f(X) = c_0 + c_1X + c_2X^2 + \dots + c_{n-1}X^{n-1}$  be the unique polynomial of degree  $\leq n - 1$  realizing  $\xi$ . Since  $(0, y_2, y_4, \dots, y_{2n-2}) = (0, a, ax_2, \dots, ax_{n-1})$  is a cycle for  $(f \circ f)(X)$ , by Remark 7 we get  $a \mid c_2, a^2 \mid c_3, \dots, a^{n-2} \mid c_{n-1}$ . If at least one number from  $c_2, c_3, \dots, c_{n-1}$  is nonzero we are done by Lemma 3.

**Second way.** We have  $y_4 - y_1 = y_4 - 1 \mid y_6 - y_0 = y_6$ , and equivalently  $ax_2 - 1 \mid ax_3$ . Hence  $ax_2 - 1 \mid ax_3x_2$  and  $ax_2 - 1 \mid x_3$  follows. If  $n > 3$ , then  $x_2 \neq 0, 1; x_3 \neq 0$ , and we are done by Lemma 3.

For  $n \geq 4$  the set  $\mathcal{Y}$  may be established by the second way.

Let  $n = 3$ . If the Lagrange interpolation polynomial  $f(X)$  realizing  $\xi = (0, 1, x_2)$  is of degree 2, then we establish  $\mathcal{Y}$  using the first way. Assume that  $f(X) = c_0 + c_1X$  realizes the cycle  $\xi$ . Then  $c_0 = 1$  and (since  $f^{\circ 3}(0) = 0$ )  $1 + c_1 + c_1^2 = 0$ . Thus  $c_1$  is a primitive third root of unity.

It remains therefore to consider  $\xi$  of the form  $\xi = (0, 1, 1 + \zeta)$ , where in what follows  $\zeta$  is a primitive third root of unity. Let  $\eta = (0, 1, y_2, \dots, y_5)$  be a cycle connected to  $(0, 1, 1 + \zeta)$ . We may write  $\eta$  in the form  $(0, 1, a, 1 + b, a(1 + \zeta), 1 + bz)$  for some  $a, b, z \in Z_K$ . Let  $g(X)$  be a polynomial realizing the cycle  $\eta$ .

We see that  $(0, 1, z)$  is a cycle for  $1/b (g^{\circ 2}(bX + 1) - 1) \in Z_K[X]$ . Thus  $(0, 1, z) \in \mathcal{S}(3)$ .

Fix  $z \neq 1 + \zeta, 1 + \zeta^2$  such that  $(0, 1, z) \in \mathcal{S}(3)$ . Owing to the invertibility of  $a-1$  we obtain that  $\eta_1 = (0, 1, b/(a-1), (a(1+\zeta)-1)/(a-1), bz/(a-1), -1/(a-1))$  is a cycle (for  $1/(a-1) (g((a-1)X + 1) - 1) \in Z_K[X]$ ) connected to  $(0, 1, z)$ . Since  $z - 1$  is not a primitive third root of unity, by the previous part of the proof we conclude that  $\eta_1$  can be effectively found, and therefore also  $a, b$  can be effectively found.

It remains therefore to consider the cases  $z = 1 + \zeta$  and  $z = 1 + \zeta^2$ .

First, let  $z = 1 + \zeta^2$ , and consider  $(0, 1, a, 1 + b, a(1 + \zeta), 1 + b(1 + \zeta^2)) \in \mathcal{S}(6)$ . Since  $y_4 - 1 \mid y_3, y_5 - y_2$  we have  $a(1 + \zeta) - 1 \mid 1 + b, 1 + b(1 + \zeta^2) - a$ . This easily gives  $a(1 + \zeta) - 1 \mid 2 + \zeta^2 \neq 0$ , which together with  $a - 1 \mid 1$  shows by Lemma 2 that  $a$  belongs to an effectively computable and finite set.

Secondly, let  $z = 1 + \zeta$ , and consider  $(0, 1, a, 1 + b, a(1 + \zeta), 1 + b(1 + \zeta)) \in \mathcal{S}(6)$  with  $a \sim b$ . Since  $y_2 - y_3$  and  $y_4 - y_5$  are units, we get that  $a - 1 - b$  and  $a(1 + \zeta) - (1 + b(1 + \zeta))$  are units. This gives that  $(a - 1 - b, a(1 + \zeta) - (1 + b(1 + \zeta)))$  lies in the finite and effectively computable set  $\mathcal{T}_K(1, \zeta; \zeta^2) = \{(u_1, v_1), \dots, (u_t, v_t)\}$ . Put  $a - 1 - b = u_i$  for some  $i \leq t$ .

If  $u_i \neq -1$ , then  $a \mid a - b = u_i + 1$ , which together with  $a - 1 \mid 1$  shows by Lemma 2 that  $a$  belongs to an effectively computable and finite set.

Let  $u_i = -1$ , or equivalently  $a = b$ , and consider  $(0, 1, a, 1 + a, a(1 + \zeta), 1 + a(1 + \zeta)) \in \mathcal{S}(6)$ . Then  $a - 1$  and  $1 + a(1 + \zeta)$  are units, and  $(a - 1, 1 + a(1 + \zeta))$  belongs to the finite and effectively computable set  $\mathcal{T}_K(1, \zeta; \zeta - 1)$ .

In this way we showed for any  $n \geq 3$  that having at our disposal  $\mathcal{S}(n)$  we can effectively find  $\mathcal{S}(2n)$ . It remains to find effectively  $\mathcal{S}(4)$ , and this case requires a slightly different approach.

Let us arbitrary order the units of the form  $\zeta_M^{i_0} \eta_1^{i_1} \cdot \dots \cdot \eta_{r+s-1}^{i_{r+s-1}}$ , with  $0 \leq i_0, i_1, \dots, i_{r+s-1} \leq 1$  and denote them as  $\sigma_1, \sigma_2, \dots, \sigma_{2^{r+s}}$ . Any unit  $\delta$  may be uniquely written in the form  $\sigma_i \epsilon^2$  for a unit  $\epsilon$  and some  $i \in [1, 2^{r+s}]$ .

Let  $(0, 1, x_2, x_3)$  be a cycle. Since  $x_2 - 1$  and  $x_3$  are units, we may write  $x_2 = 1 + \delta, x_3 = \epsilon$  for some invertible  $\delta, \epsilon$ . Thus  $(0, 1, x_2, x_3) = (0, 1, 1 + \delta, \epsilon)$ . Since  $x_2 \sim 1 - x_3$ , we may write  $1 - \epsilon = \psi(1 + \delta)$  for some unit  $\psi$ . Since  $x_2 - x_3$  is invertible,  $(x_2 - x_3)/\psi = 1 + \delta + \delta/\psi$  is invertible as well. In view of  $\delta + \delta^2 - \delta/\psi = -\delta\epsilon/\psi$ , we then obtain that

$$\tau_1 := 1 + \delta + \frac{\delta}{\psi}; \quad \tau_2 := \delta + \delta^2 - \frac{\delta}{\psi}$$

are units. Write  $\tau_1 = \sigma_i \rho^2$ , for some  $1 \leq i \leq 2^{r+s}$  and a unit  $\rho$ .

We see that  $\tau_2 = (\delta + 1)^2 - \tau_1 = (\delta + 1)^2 - \sigma_i \rho^2 = (\delta + 1 + \sqrt{\sigma_i} \rho)(\delta + 1 - \sqrt{\sigma_i} \rho)$  is invertible. Thus  $\delta + 1 + \sqrt{\sigma_i} \rho$  and  $\delta + 1 - \sqrt{\sigma_i} \rho$  are units in  $Z_{K(\sqrt{\sigma_i})}$ . Hence

$$\left( \frac{\delta + 1}{\sqrt{\sigma_i} \rho} + 1, \frac{\delta + 1}{\sqrt{\sigma_i} \rho} - 1 \right) \in \mathcal{T}_{K(\sqrt{\sigma_i})}(1, -1; 2).$$

For  $i = 1, \dots, 2^{r+s}$  the sets  $\mathcal{T}_{K(\sqrt{\sigma_i})}(1, -1; 2)$  are finite and effectively computable. Moreover, having  $K$  we may effectively find all  $K(\sqrt{\sigma_i})$ .

Let  $\mathcal{T}_{K(\sqrt{\sigma_i})}(1, -1; 2) = \{(u_{i_1}, v_{i_1}), \dots, (u_{i_{j(i)}}, v_{i_{j(i)}})\}$ . Hence  $(\delta + 1)/(\sqrt{\sigma_i} \rho) + 1 = u_{i_j}$  for some  $j \leq j(i)$ . Thus  $(u_{i_j} - 1)\sqrt{\sigma_i} \in K$ , and therefore  $(\delta, \rho) \in \mathcal{T}_K(1, \sqrt{\sigma_i}(1 - u_{i_j}); -1)$ .

For any  $i \leq 2^{r+s}$  and  $u_{i_j}$  satisfying  $(u_{i_j} - 1)\sqrt{\sigma_i} \in K$  we effectively find  $\mathcal{T}_K(1, \sqrt{\sigma_i}(1 - u_{i_j}); -1)$ , and this gives that all possible  $x_2 - 1 = \delta$  belong to a finite and effectively computable set.

Having found possibilities for  $\delta$  and observing that  $(\tau_1, -\delta/\psi) \in \mathcal{T}_K(1, 1; 1 + \delta)$ , we finally obtain that  $x_3 = \epsilon = 1 - \psi(1 + \delta)$  also belongs to some finite and effectively computable set. The proof of Theorem 2 is thus completed.  $\square$

### 3. Finite orbits

In a domain  $R$  a tuple  $(y_k, y_{k-1}, \dots, y_0 = x_0, x_1, \dots, x_{n-1})$  of distinct elements from  $R$  is called a (*finite*) *orbit* provided there exists a polynomial  $f(X) \in R[X]$  realizing this orbit, i.e.  $f(y_k) = y_{k-1}$ ,  $f(y_{k-1}) = y_{k-2}, \dots, f(y_1) = y_0 = x_0$ ,  $f(x_0) = x_1$ ,  $f(x_1) = x_2, \dots, f(x_{n-1}) = x_0$ . We underlined the unique cycle contained in the orbit.

The counterpart of the second part of Remark 1 holds also for orbits.

The number  $n + k$  is called the *length* of this orbit, the cycle  $(x_0, x_1, \dots, x_{n-1})$  will be called the *head* of this orbit (of length  $n$ ),  $(y_k, y_{k-1}, \dots, y_0)$  will be called the *tail* of this orbit (of length  $k$  (not  $k + 1$ )), and finally  $(n, k)$  will be called the *type* of this orbit.

A cycle  $(x_0, x_1, \dots, x_{n-1})$  in  $R$  is called *linear*, provided it is realized by some polynomial  $f(X) \in R[X]$  of degree  $\leq 1$ . We call an orbit *linear* provided its head is linear.

If  $(y_k, \dots, y_0 = x_0, \dots, x_{n-1})$  is an orbit in  $R$ , then clearly for any invertible  $a \in R$  and any  $b \in R$  the tuple  $(ay_k + b, \dots, ay_0 + b = ax_0 + b, \dots, ax_{n-1} + b)$  is also an orbit in  $R$ . Two such orbits will be called *equivalent*.

In [H-KNa2] it was established that in any finitely generated domain of characteristic zero there is only finitely many inequivalent nonlinear orbits.

**Theorem 3.** *Let  $K$  be an algebraic number field of degree  $N$ . Then for polynomial orbits in  $Z_K$  the following holds.*

- (i) *The lengths of orbits are bounded by some quantity depending solely on  $N$ .*
- (ii) *There are only finitely many inequivalent nonlinear orbits and all of them can be effectively found.*
- (iii) *Any linear orbit with tail of length 0 is equivalent to  $(0, a, a(1 + \zeta_n), \dots, a(1 + \zeta_n + \zeta_n^2 + \dots + \zeta_n^{n-2}))$  for some nonzero  $a \in Z_K$  and some primitive  $n$ -th root of unity  $\zeta_n \in Z_K$ .*
- (iv) *There are only finitely many inequivalent linear orbits with head of length  $\geq 4$  and tail of length  $\geq 1$  and all of them can be effectively found.*
- (v) *Any linear orbit of type  $(3, 1)$  is equivalent to  $(1, 0, 1 + \epsilon, (1 + \epsilon)(1 + \zeta_3))$  (this is the orbit for  $f(X) = (X - 1)(X - (1 + \epsilon)(1 + \zeta_3))(-\zeta_3 + (\zeta_3/\epsilon)X)$ , and  $\zeta_3$  is a primitive third root of unity) for any unit  $\epsilon \neq -1$ .*
- (vi) *There are only finitely many inequivalent linear orbits of type  $(3, k)$  with  $k \geq 2$ , and all of them can be effectively found.*
- (vii) *Any orbit of type  $(2, 1)$  is equivalent to  $(1, 0, d)$  for some  $d \in Z_K, d \neq 0, 1$  (this is the orbit for  $f(X) = (d - X)(1 - X)$ ).*
- (viii) *Any orbit of type  $(2, 2)$  is equivalent to  $(1 + \epsilon, 1, 0, 1 + \epsilon + \delta)$  for some invertible  $\epsilon, \delta \in Z_K, \epsilon \neq -1, \delta \neq -\epsilon, -1 - \epsilon; (1 + \epsilon) \mid \delta - 1$  (this is the orbit for  $f(X) = (1 + \epsilon + \delta - X)(1 - X) - (1 + \epsilon\delta)/(\epsilon\delta(1 + \epsilon))X(X - 1)(X - (1 + \epsilon + \delta))$ ).*
- (ix)  *$(a, b)$  is the orbit of type  $(1, 1)$  for any  $a \neq b$  (for  $f(X) = b$ ).*
- (x) *Any orbit of type  $(1, 2)$  is equivalent to  $(d, d(1 - \epsilon), 0)$  for some nonzero  $d \in Z_K$  and invertible  $\epsilon$  satisfying  $d \mid 1 - \epsilon$  (this is the orbit for  $f(X) = (1 - \epsilon)/(\epsilon d)(X - d(1 - \epsilon))X$ ).*
- (xi) *There are only finitely many inequivalent orbits with head of length 1 or 2 and tail of length  $\geq 3$ . One can effectively (up to equivalence) find all of them if and only if one finds an effective procedure for determining all solutions of  $u + v + w = 1, 1 - u \mid 1 - v$ , with invertible  $u, v, w \in Z_K, u, v, w \neq 1$ .*

PROOF. Any orbit is equivalent to  $(y_k, \dots, y_1, y_0 = 0, a, ax_2, \dots, ax_{n-1})$ , where  $(0, 1, x_2, \dots, x_{n-1})$  is a normalized cycle, and we may restrict considerations to orbits of this form.

(i) In [NaPe] it was emphasized that the lengths of orbits in  $Z_K$  are bounded by some quantity depending solely on  $B(Z_K)$  and the number of nontrivial solutions of  $u + v + w = 1$  in units (the latter number is bounded by some expression depending solely on  $[K : Q]$  (see [EG])).

**Lemma 4.** Fix  $a, b \in Z_K$ ,  $a \neq 0$ ,  $b \neq 0$ ,  $a \neq b$ .

- ( $\alpha$ ) There are only finitely many orbits with head of length  $\geq 3$  of the form  $(y_k, \dots, \underline{y_0 = 0, a, \dots})$ , and all of them can be effectively found.
- ( $\beta$ ) There are only finitely many orbits of the form  $(y_k, \dots, y_1 = b, \underline{y_0 = 0, a})$ , and all of them can be effectively found.
- ( $\gamma$ ) There are only finitely many orbits of the form  $(y_k, \dots, y_1 = a, \underline{y_0 = 0})$ , and all of them can be effectively found.

PROOF. ( $\alpha$ ) Put  $(x_0, x_1, \dots, x_{n-1}) = (0, a, \dots)$ . Then  $(0, 1, x_2/x_1, x_3/x_1, \dots)$  is a normalized cycle. By Theorem 2 there is only finitely many possibilities for  $n \geq 3$ , then for  $x_2, \dots, x_{n-1}$  and they can be effectively computed. It suffices then to deal with tails, which have bounded lengths by (i) of Theorem 3.

Since for  $x_2$  there is only finitely many possibilities, we may fix  $x_2 = c$ . The assertion follows then from Theorem 4 of [H-KNa2], as all solutions of the unit equations  $a_1u + b_1v = c_1$  (with nonzero  $a_1, b_1, c_1$ ) can be found in an effective way. One may also use Lemma 2.

( $\beta$ ) and ( $\gamma$ ) follow immediately from Theorem 3(i) and Theorem 4 of [H-KNa2].  $\square$

(ii) Since we are considering orbits up to equivalence, by Theorem 2, Remark 7 and (C) we may assume that the element  $a$  from a nonlinear orbit  $(y_k, \dots, y_1, \underline{y_0 = 0, a, ax_2, \dots, ax_{n-1}})$  belongs to some effectively computable and finite set  $\mathcal{Y}$ . By Lemma 4( $\alpha$ ) we are done.

(iii) It is clear.

(iv), (v) Let  $(b, 0, a, a(1 + \zeta_n), \dots, a(1 + \zeta_n + \dots + \zeta_n^{n-2}))$  be a linear orbit (for a polynomial  $f(X) \in Z_K[X]$ ) with head of length  $n \geq 3$ , and  $\zeta_n$  is a primitive  $n$ -th root of unity. Thus  $f(X)$  is of the form  $f(X) = \zeta_n X + a + h(X)X(X - a) \cdot \dots \cdot (X - a(1 + \zeta_n + \dots + \zeta_n^{n-2}))$  for some  $h(X) \in Z_K[X]$ . Since  $b - 0 \mid f(b) - f(0) = 0 - a$  we may write  $a = bd$ , for some  $d \in Z_K$ . In view of  $h(b) \in Z_K$  we get  $b^{n-1}(1 - d)(1 - d(1 + \zeta_n)) \cdot \dots \cdot (1 - d(1 + \zeta_n + \dots + \zeta_n^{n-2})) \mid \zeta_n + d \neq 0$  (otherwise  $y_1 = x_{n-1}$ ).

For  $n \geq 4$  we then get  $1 - d \mid d + \zeta_n$  and  $(1 - d(1 + \zeta_n)) \mid d + \zeta_n$ , which gives  $1 - d \mid \zeta_n + 1$  and  $(1 - d(1 + \zeta_n)) \mid \zeta_n(1 + \zeta_n) + 1 \neq 0$ . Lemma 2 gives  $d \in \mathcal{Y}$  for some finite and effectively computable  $\mathcal{Y}$ . Fix any  $d \in \mathcal{Y}$ , and we obtain  $b \mid b^{n-1} \mid \zeta_n + d$ . By (C),  $b$  is associated to an element of some finite and effectively computable  $\mathcal{Y}_1$ . Thus our orbit is equivalent to some orbit of the form  $(b_1, 0, b_1 d, \dots)$  with  $b_1 \in \mathcal{Y}_1, d \in \mathcal{Y}$ . Since there is only finitely many possibilities for  $b_1 d$ , (iv) follows from Lemma 4( $\alpha$ ).

Let  $n = 3$ . Then  $b^2(1-d)(1-d(1+\zeta_3)) \mid d + \zeta_3$ , which gives that  $b^2(1-d)$  is invertible. So our orbit is equivalent to  $(1, \underline{0}, \underline{d}, \underline{d(1+\zeta_3)})$  with invertible  $d-1$ . This settles (v).

(vi) Consider a linear orbit  $(y_k, \dots, y_1, \underline{0}, \underline{x_1}, \underline{x_2})$  of type  $(3, k)$  with some  $k \geq 2$ . By (v) we may assume that  $y_1 = 1$ ,  $x_1 = d$ ,  $x_2 = d(1+\zeta_3)$ , with  $d = 1 + \epsilon$  for some invertible  $\epsilon \neq -1$ . So  $(y_2, \underline{1}, \underline{0}, \underline{d}, \underline{d(1+\zeta_3)})$  is an orbit. This gives  $y_2 - 1 \mid 1 - 0 = 1$ ;  $y_2 = y_2 - 0 \mid 1 - (1 + \epsilon) \mid 1$ . By Lemma 2, we obtain  $y_2 \in \mathcal{Y}$ , for some finite and effectively computable  $\mathcal{Y}$ .

Fix  $y_2 \in \mathcal{Y} \setminus \{-\zeta_3\}$ . Then  $y_2 - d \mid d(1+\zeta_3) - 1$ , and  $y_2 - d \mid \zeta_3 + y_2$  follows. This together with  $d-1 \mid 1$ , by Lemma 2, gives that  $d$  belongs to some finite and effectively computable set. Lemma 4( $\alpha$ ) gives the assertion for such  $y_2$ .

Assume now that  $y_2 = -\zeta_3$ . Then  $d(1+\zeta_3) - (-\zeta_3) \mid 1 - 0 = 1$ . This together with  $d-1 \mid 1$ , by Lemma 2, gives that  $d$  belongs to some finite and effectively computable set. Lemma 4( $\alpha$ ) gives the assertion for  $y_2 = -\zeta_3$ . This settles (vi).

(vii) Let  $(c, \underline{0}, \underline{b})$  be an orbit for  $f(X)$ . So  $f(X) = b - X + h(X)X(X - b)$  for some  $h(X) \in Z_K[X]$ . This gives  $c(c-b) \mid b-c$ , and  $c$  is invertible. We may thus assume that  $c = 1$ .

(viii) By (vii) any orbit of type  $(2, 2)$  is equivalent to  $(m, \underline{1}, \underline{0}, \underline{d})$ . This gives  $m-1 \mid 1$ ;  $m-d \mid 1$  and  $m \mid d-1$ .

(ix) It is obvious.

(x) Let  $(d, c, \underline{0})$  be an orbit for  $f(X)$ . So  $f(X) = X(X-c)h(X)$  for some  $h(X) \in Z_K[X]$ . This gives  $d(d-c) \mid c$ , and  $d \sim d-c$  follows. Put  $d-c = d\epsilon$ . The rest is obvious.

(xi) Let us first deal with orbits with head of length 1.

*Suppose, for the time being, that we have at our disposal all orbits of the form  $(1, a, b, \underline{0})$ , and there is only finitely many of them. (\*)*

Let  $(m, c, d, \underline{0})$  be an orbit for some  $f(X)$ . We see that  $(1, c/m, d/m, \underline{0})$  is the orbit for  $(1/m)f(mX) \in Z_K[X]$ . Thus  $(1, c/m, d/m, \underline{0}) = (1, a, b, \underline{0})$ , with specified possible values for  $a, b$ . Since  $(c, d, \underline{0})$  is the orbit, by (x), we have  $m \mid c \mid (d/c) = (b/a)$ . This gives  $d = m(d/m) \mid (b/a)b = (b^2)/a$ . Thus in any orbit of the form  $(y_k, \dots, y_1, \underline{0})$  (with  $k \geq 3$ ) the number  $y_1$  (still under our assumption  $(*)$ ) may assume (up to associates) only finitely many known values. By Lemma 4( $\gamma$ ) we then would be able to find effectively (up to equivalence) all orbits with head of length 1 and tail of length  $\geq 3$ .

So let  $(1, a, b, \underline{0})$  be an orbit. We easily obtain  $1-a \mid a$ ;  $a \mid b$ ;  $a-b \mid b$ ;  $1-b \mid b$ . This gives  $a = 1 - \delta$ ,  $b = a(1 - \epsilon)$ ,  $b = 1 - \psi$  for some invertible  $\delta, \epsilon, \psi \neq 1$ , and by (x) we get  $1 - \delta \mid 1 - \epsilon$ . This gives  $\psi = \delta + \epsilon - \delta\epsilon$  and

$1/\delta + 1/\epsilon - \psi/(\delta\epsilon) = 1$  follows. Hence  $(1/\delta, 1/\epsilon, -\psi/(\delta\epsilon))$  is the solution of the 3-unit equation  $u + v + w = 1$  satisfying  $(1 - 1/\delta) \mid (1 - 1/\epsilon)$ .

If this solution is trivial, then  $\delta = -\epsilon$  and we obtain  $1 - \delta \mid 1 + \delta$ , i.e.  $1 - \delta \mid 2$ . By Lemma 2, we may effectively find all such  $\delta$ , and there is only finitely many of them.

Conversely, suppose that  $(u, v, w)$  is a nontrivial solution of the 3-unit equation  $u + v + w = 1$  with  $u - 1 \mid v - 1$ . Then  $(1, 1 - 1/u, (1 - 1/u)(1 - 1/v), \underline{0})$  is the orbit for

$$X \left( X - \left( 1 - \frac{1}{u} \right) \left( 1 - \frac{1}{v} \right) \right) \times \left( \frac{u(v-1)}{u-1} + \frac{u(v-uv^2+u^2-u)}{(u+v-1)(u-1)} \left( X - \left( 1 - \frac{1}{u} \right) \right) \right) \in Z_K[X].$$

Notice that if  $(u_1, v_1, w_1) \neq (u_2, v_2, w_2)$ , then  $(1, 1 - 1/u_1, (1 - 1/u_1)(1 - 1/v_1), \underline{0})$  is not equivalent to  $(1, 1 - 1/u_2, (1 - 1/u_2)(1 - 1/v_2), \underline{0})$ .

In this way we obtained the one-to-one correspondence between the set of all nontrivial solutions of the 3-unit equation  $u + v + w = 1$  satisfying  $u - 1 \mid v - 1$  and a certain subset  $\mathcal{S}$  of orbits of the form  $(1, a, b, \underline{0})$ . Since any orbit of the form  $(1, a, b, \underline{0})$  lying out of  $\mathcal{S}$  is effectively computable we are done in the case of orbits with head of length 1.

Now we deal with orbits with head of length 2.

By (viii) and Lemma 4( $\beta$ ) it suffices to find all orbits of the form  $(t, m, 1, \underline{0}, d)$  provided we have all nontrivial solutions of the 3-unit equation  $u + v + w = 1$  with  $u - 1 \mid v - 1$ .

*Assume that we have all such solutions of this 3-unit equation at our disposal. (\*\*)*

Since  $t - m \mid m - 1 \mid 1 - 0 = 1$ ;  $t - d \mid m - 0 = m$ ;  $t - 1 \mid m - 0 = m$ ;  $t = t - 0 \mid m - d \mid 1 - 0 = 1$ , we obtain that  $t, t - m, m - 1, m - d$  are invertible and  $1 - t \mid m$ . Hence  $(t, 1 - m, m - t)$  is the solution of the 3-unit equation  $u + v + w = 1$  satisfying  $t - 1 \mid (1 - m) - 1 = -m$ .

If this solution is nontrivial, then by (\*\*) the numbers  $t, m$  belong to some finite and explicitly given set.

If this solution is trivial, then  $t = m - 1$ . This gives  $m - 2 \mid m$ , i.e.  $m - 2 \mid 2$  and  $m - 1 \mid 1$ . By Lemma 2 we obtain that for such  $m$  there is only finitely many possibilities, and they can be effectively computed.

All in all, there is only finitely many possibilities for  $t, m$ , and all these possibilities can be effectively computed assuming (\*\*). Fix any possible  $t, m$ .

Then  $t - d \mid m$ ;  $m - d \mid 1$  and by Lemma 2 we may compute all possible values for  $d$ . This settles (xi).  $\square$

*Remark 9.* In the proof of Theorem 3(xi) we showed that in terms of effective computability finding (up to equivalence) all orbits of type  $(1, k)$  (with  $k \geq 3$ ) in an effective way is equivalent to finding in an effective way all nontrivial solutions of  $u + v + w = 1$  satisfying  $u - 1 \mid v - 1$ . One sees that finding in an effective way all orbits (up to equivalence) of type  $(2, k)$  (with  $k \geq 3$ ) is equivalent to finding in an effective way all nontrivial solutions of  $u + v + w = 1$  satisfying  $u - 1 \mid v - 1$  and some other condition.

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