

## Polynomial values in linear recurrences

I. NEMES and A. PETHŐ (Debrecen)

### 1. Introduction

Let  $A_1, \dots, A_k$  and  $G_0, G_1, \dots, G_{k-1}$  be integers. We have for the  $n$ -th term of a  $k$ -order linear recurrence

$$(1) \quad G_n = A_1 G_{n-1} + \dots + A_k G_{n-k} \quad \text{for } n = k, k+1, \dots$$

Let  $\alpha_1, \dots, \alpha_t$  be the distinct roots of the characteristic polynomial of the recurrence

$$(2) \quad X^k - A_1 X^{k-1} - \dots - A_k.$$

Throughout this paper we assume that  $\alpha_1$  has multiplicity one. Then for  $n \geq 0$  we have

$$(3) \quad G_n = E_1 \alpha_1^n + P_2(n) \alpha_2^n + \dots + P_t(n) \alpha_t^n,$$

where  $P_i(n)$  is a polynomial with degree less than the multiplicity of  $\alpha_i$  in the characteristic polynomial of  $G_n$ , and where  $E_1$  and the coefficients of  $P_i(n)$  are elements of the field  $\mathbb{Q}(\alpha_1, \dots, \alpha_t)$ .

Finally let  $T(x) = B_m x^m + \dots + B_0$  be a polynomial with integer coefficients. Its degree will be denoted by  $\deg T$ , while its height,  $\max \{|B_i|; i=0, \dots, m\}$  by  $\mathcal{H}(T)$ .

The Diophantine equation

$$(4) \quad G_n = E x^q + T(x) E \neq 0, \quad \text{integer}$$

was investigated by several authors. Naturally, most of the results are known for  $T(x) = 0$ .

SHOREY and STEWART [4] proved for general linear recurrences that (4) has finitely many solutions in  $q$ , assuming  $|\alpha_1| > |\alpha_j|$ ,  $j=2, \dots, t$ . Under some other restriction on  $G_n$ , recently P. KISS [2] was able to generalize their result when  $\deg T < c_1 q$ .

For nondegenerate second order linear recurrences SHOREY and STEWART [4] derived much more, namely (4) has finitely many solutions in integers  $|x| > 1$ ,  $q \geq 2, n$ . The second author investigated in [3] for nondegenerate second order linear recurrences the slightly more general equation

$$(5) \quad G_n = w x^q$$

with  $w \in S$ , where  $S$  is the set of nonzero integers composed solely of primes from some fixed finite set. He showed assuming yet  $(A_1, A_2) = 1$  that (5) has finitely many solutions in integers  $|x| > 1$ ,  $q \geq 2$ ,  $n$  and  $w \in S$ .

STEWART [7] was dealing with the case  $T(x) = c$ , where  $c$  is a fixed integer. For nondegenerate second order recurrences with  $|A_2| = 1$  he proved the finiteness of the solution in integers  $|x| > 1$ ,  $q > 2$ ,  $c, n$  of (4).

All the above mentioned results are effective.

In the present paper we shall derive some results for (4), when  $T(x)$  is a polynomial with some restriction.

## 2. Main results

**Theorem 1.** Let  $G_n$  satisfy (3),  $\alpha_1, \alpha_2 \neq 1$ ,  $|\alpha_1| > |\alpha_2| > |\alpha_j|$ ,  $j = 3, \dots, t$  and  $G_n - E_1 \alpha_1^n \neq 0$  for  $n > c_2$ . Further let  $\mathcal{H}(T) < H_1$  and  $\deg T \leq qc_3$ , where  $H_1 > 0$  real number. Then all integer solutions  $n, |x| > 1$ ,  $q \geq 2$  of the equation (4) satisfy  $q < c_4$ , where  $c_2, c_3$  and  $c_4$  are effectively computable constants depending on  $E, G_n$  and  $H_1$ .

For second order linear recurrences we prove a more precise result.

**Theorem 2.** Let  $G_n$  be a nondegenerate second order linear recurrence with  $|A_2| = 1$ . Further let  $\mathcal{H}(T) < H_2$  and  $\deg T \leq \min\{q(1-\gamma), q-3\}$  where  $H_2$  and  $\gamma < 1$  are positive real numbers. Then all integer solutions  $n, |x| > 1$ ,  $q \geq 2$  of (4) satisfy

$$\max\{n, |x|, q\} < c_5,$$

where  $c_5$  is an effectively computable constant depending on  $E, G_n, \gamma$  and  $H_2$ .

*Remark.* Theorem 2 is in the restriction of  $\deg T$  best possible. Let  $L_n$  denote a Lucas sequence, i.e.  $L_0 = 2$ ,  $L_1 = \alpha + \beta$  and  $L_n = (\alpha + \beta)L_{n-1} + L_{n-2}$ , where  $\alpha\beta = -1$  and  $\alpha + \beta$  integer. Then, as is well known,  $L_n = \alpha^n + \beta^n$ . Further it is easy to see that  $L_{2n} = L_n^2 + (-1)^n 2$ . This means that both equations  $L_n = x^2 + 2$  and  $L_n = x^2 - 2$  have infinitely many integer solutions  $n, x$ . Therefore in Theorem 2 the assumption  $\deg T \leq q - 3$  is necessary.

## 3. Auxiliary results

The most important result we use is Lemma 6 of [4].

**Lemma.** Let  $\alpha$  be a real algebraic number larger than one from the field  $k$ . Let  $[K:Q] = D, E, A$  and  $B$  be elements of  $K, EAB \neq 0$ , finally  $\delta$  a positive real number. If  $Ex^q = A\alpha^n + B$  with  $|B| < \alpha^{n(1-\delta)}$  and  $n, x, q$  integers larger than one then  $q < c_6$  is a constant, effectively computable in terms of  $D, E, A, \alpha$  and  $\delta$ .

The following theorem was proved by C. L. SIEGEL [5] for the first time but in noneffective form. Using the upper estimate for linear forms of logarithms of algebraic numbers A. BAKER [1] proved it in effective form.

**Theorem A.** *Let  $F(x)$  be a monic polynomial of degree  $n$  and with integer coefficients. Let it have at least three simple zeros. Then the integer solutions  $x, y$  of the equation  $Ay^2 = F(x)$  where  $A$  is an integer, satisfy  $\max\{|x|, |y|\} < c_7$  a constant effectively computable in terms of  $A$  and the coefficients of  $F(x)$ .*

#### 4. Proofs

In the sequel  $c_8, c_9, \dots$  will denote positive numbers effectively computable in terms of  $E, G_n, H_i$  and  $\gamma$ .

**PROOF OF THEOREM 1.** We may assume  $\alpha_1$  to be positive by changing if necessary the sign of  $E_1$ . Further since  $\alpha_1$  is an algebraic integer with absolute value strictly larger than all its conjugates, on taking norm we see that either  $\alpha_1 > 1$  or  $\alpha_1$  is one of 0 or 1. But these two cases were excluded so we may assume  $\alpha_1 > 1$ . Put

$$B_1(n) = P_2(n)\alpha_2^n + \dots + P_t(n)\alpha_t^n$$

$$D_1 = \max\{\deg P_i(n); i = 2, \dots, t\}.$$

It is easy to show that

$$(6) \quad 2|B_1| < c_8 n^{D_1} |\alpha_2|^n.$$

Assume that for a polynomial  $T(x)$  with  $\mathcal{H}(T) \leq H_1$  and  $\deg T < qc_3, n, |x| > 1, q$  is a solution of (4). We shall give an estimate for  $c_3$  in the proof.

Write (4) in the form

$$Ex^q = E_1\alpha_1^n + B_1(n) - T(x).$$

Assume first that

$$(7) \quad B_1(n) - T(x) = 0$$

in which case

$$(8) \quad E_1\alpha_1^n = Ex^q$$

also holds. We distinguish two cases.

If  $1 > |\alpha_2| \cong |\alpha_j|, j=3, \dots, t$ , then  $\lim_{n \rightarrow \infty} B_1(n) = 0$  which means  $|B_1(n)| < 1$  for  $n > c_8$ . Further  $T(x)$  is a polynomial with integer coefficients, therefore (7) has for  $n > c_8$  no solutions. By (8)  $q = c_9 \frac{n}{\log |x|} < c_{10}$  because of  $|x| \cong 2$ .

If  $|\alpha_2| > 1$ , then write

$$B_1(n) = P_2(n)\alpha_2^n \left( 1 + \sum_{i=3}^t \frac{P_i(n)}{P_2(n)} \left( \frac{\alpha_i}{\alpha_2} \right)^n \right).$$

The quantity in the brackets tends to 1 if  $n$  tends to infinity, so for  $n > c_{11}$   $|B_1(n)| > |P_2(n)||\alpha_2|^n(1 - \varepsilon') > |\alpha_2|^{n(1-\varepsilon)}$ . On the other hand  $T(x) \leq m H_1 |x|^m$ . From (8) we have  $|x| = (|E_1/E||\alpha_1|^n)^{1/q}$ , hence  $|T(x)| \leq c_{12} m |\alpha_1|^{nm/q} < |\alpha_1|^{c_{13}nm/q}$ . By (7)  $|\alpha_2|^{n(1-\varepsilon)} < |B_1(n)| = |T(x)| < |\alpha_1|^{c_{13}nm/q}$ . This implies  $m > c_{14} \frac{\log |\alpha_2|}{\log |\alpha_1|} q$ . Therefore

if  $c_3 \leq c_{13} \frac{\log |\alpha_2|}{\log |\alpha_1|}$  then (7) and (8) have only finitely many solutions in  $n, q, |x| > 1$  which are effectively computable.

In the sequel we assume that  $B_1(n) - T(x) \neq 0$ . Put  $\delta = (1 - \vartheta)/2$ , where  $\vartheta = 0$  if  $\alpha_2 \equiv 1$ , and  $\vartheta = \frac{\log |\alpha_2|}{\log |\alpha_1|}$  otherwise. Since  $\alpha_1 > 1$  we have by (6)  $|B_1| < \frac{1}{2} \alpha_1^{n(1-\delta)}$  for  $n > c_{14}$ . We choose  $c_{15}$  such that  $|E||x|^{c_{15}} > 2E_1 + 1$ . If  $q$  is large enough then  $\frac{|T(x)|}{|x|^{q-c_{15}}} < 1$ . Write (4) in the form

$$Ex^{c_{15}} = \frac{E_1 \alpha_1^n + B_1(n)}{x^{q-c_{15}}} - \frac{T(x)}{x^{q-c_{15}}}.$$

Taking absolute values, and applying the above estimates we have

$$2|E_1| + 1 < |Ex^{c_{15}}| \leq \left| \frac{E_1 \alpha_1^n + B_1(n)}{x^{q-c_{15}}} \right| + 1.$$

Hence  $2|E_1||x|^{q-c_{15}} \leq |E_1 \alpha_1^n + B_1(n)| \leq 2|E_1| \alpha_1^n$ , so  $|x|^{q-c_{15}} < \alpha_1^n$ . Further  $|T(x)| \leq m H_1 |x|^m < |x|^{c_{16}q}$ , with  $c_{16} < 1$ . Hence  $|T(x)| < \alpha_1^{c_{16}qn/(q-c_{15})} < \frac{1}{2} \alpha_1^{n(1-\delta)}$  when  $\delta$  is small enough. So we have  $|B_1(n) + T(x)| \leq \alpha_1^{n(1-\delta)}$ .

Note that if  $n < c_{17}$  and (4) holds then  $q < c_{18}$  as required. Of course since

$$|x|^m (Ex^{q-m} - mH_1) \leq |E||x|^q - |T(x)| \leq |Ex^q + T(x)| = |G_n| \leq c_{19} \alpha_1^n$$

the required inequality for  $q$  holds.

Finally by the Lemma if  $n > c_{20}$  and (4) holds then  $q < c_1$  as was stated.

PROOF OF THEOREM 2. The assumption  $|A_2| = 1$  means  $|\alpha_1 \alpha_2| = 1$ . We show that  $\alpha_1$  and  $\alpha_2$  are real numbers. Of course if one of them had a nonzero imaginary part then  $|\alpha_1| = |\alpha_2|$  would hold since they are roots of a polynomial with integer coefficients. This implies  $|\alpha_1| = |\bar{\alpha}_2| = 1$ .

But  $\alpha_1/\alpha_2$  cannot be a root of unity by nondegeneracy. Thus  $\alpha_1, \alpha_2$  are real numbers and  $|\alpha_1| > 1 > |\alpha_2|$  holds since  $|\alpha_1 \alpha_2| = 1$ . Further the equation  $A_2 \alpha_2^n = T(x)$  has only finitely many solutions, since with  $n$  large enough  $0 < |T(x)| = |A_2 \alpha_2^n| < 1$ . Therefore  $n < c_{21}$ , which implies  $q, |x| < c_{22}$ .

In the sequel we assume  $A_2 \alpha_2^n \neq T(x)$ . Now we shall prove  $q < c_{23}$ . Assume that (4) has a solution  $n, q, |x| > 1$  such that  $q \geq c_{23}$ , with  $\frac{\log q H_2}{q \log 2} < \frac{\gamma}{2}$ , and  $q \frac{\gamma}{4} > 1 + \frac{\gamma}{4}$ . Then

$$|T(x)| < m H_2 |x|^m < |x|^{m + \frac{\log H_2 m}{\log 2}} < |x|^{q(1-\gamma) + \frac{\log q H_2}{\log 2}}.$$

Applying the assumption we have

$$q(1-\gamma) + \frac{\log q H_2}{\log 2} < q(1-\gamma) + q \frac{\gamma}{2} = q \left(1 - \frac{\gamma}{2}\right) < q - 1.$$

Hence  $\frac{|T(x)|}{|x|^{q-1}} < 1$ . From this follows as in the proof of Theorem 1  $|x|^{q-1} < \alpha_1^n$ ,

and  $|T(x)| < \alpha_1^{(1-\delta)}$  with  $\delta = \frac{\gamma}{4}$ . Finally it is obvious that  $P_2(n)|\alpha_2|^n < \alpha_1^{n(1-\delta)}$ .

Now applying the Lemma we conclude that  $q$  is bounded by an effectively computable constant.

Now let  $q$  and  $T(x)$  be fixed with  $\deg T < q(1-\gamma)$  and consider the equation

$$G_n = Ex^q + T(x) = T_1(x).$$

It is well known that  $G_{n+1}^2 - A_1 G_{n+1} G_n + A_2 G_n^2 = CA_2^n$ , with  $C = G_1^2 - A_1 G_1 G_0 + A_2 G_0^2$  (see for example [3] Lemma 1). From this follows  $DG_n^2 + 4CA_2^n = z^2$ , with  $D = A_1^2 - 4A_2 \neq 0$  and  $z$  an integer. Replacing  $G_n^2$  by  $Ex^q + T(x)$ , and taking  $|A_2| = 1$  into account we have

$$R(x) = D(Ex^q + T(x))^2 \pm 4C = DT_1^2(x) \pm 4C = z^2.$$

This is an elliptic equation, and by means of Theorem A it has finitely many solutions in  $x, z$  when  $R(x)$  has at least three simple zeros. Let  $(R(x), R'(x)) = Q(x)$ . It is well known that a root  $\omega$  of  $R(x)$  has multiplicity at least two if and only if  $\omega$  is a root of  $Q(x)$ . Further  $(R(x), T_1(x)) = 1$  because of  $c \neq 0$ , and  $R'(x) = 2D T_1(x) T_1'(x)$  so  $\deg Q(x) \leq q-1$ . This means that if either  $\deg Q(x) < q-1$  or it has at least one multiple root then  $R(x)$  has at least three simple zeros and we are ready.

Hence the only wrong case is when  $Q(x) = T_1'(x)$ , and  $R(x) = T_1'(x)^2 S(x)$  with a polynomial  $S(x)$  with rational coefficients of degree two with not any multiple roots. Let  $S(x) = s_2 x^2 + s_1 x + s_0$  and consider the equation

$$D(Ex^q + T(x))^2 \pm 4C = (2qEx^{q-1} + T'(x))^2 (s_2 x^2 + s_1 x + s_0).$$

The coefficient of  $x^{2q-1}$  and that of  $x^{2q-2}$  on the left hand side is 0, because of  $\deg T(x) < q-3$ , while the coefficient of  $x^{2q-1}$  on the right hand side is  $4q^2 E^2 s_1$ , and that of  $x^{2q-2}$  is  $4q^2 E^2 s_0$ . This means  $s_1 = s_0 = 0$ , and  $S(x) = s_2 x^2$  which is a contradiction.

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MATHEMATICAL INSTITUT  
KOSSUTH LAJOS UNIVERSITY  
H-4010 DEBRECEN, Pf. 12.  
HUNGARY