

A note on the arithmetic properties of Stern Polynomials

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Abstract. We investigate the Stern polynomials defined by $B_0(t) = 0$, $B_1(t) = 1$, and for $n \geq 1$ by the recurrence relations $B_{2n}(t) = tB_n(t)$, $B_{2n+1}(t) = B_n(t) + B_{n+1}(t)$. We prove that all possible rational roots of that polynomials are $0, -1, -1/2, -1/3$. We give complete characterization of n such that $\deg(B_n) = \deg(B_{n+1})$ and $\deg(B_n) = \deg(B_{n+1}) = \deg(B_{n+2})$. Moreover, we present some results concerning reciprocal Stern polynomials.

1. Introduction

We consider the sequence of Stern polynomials $(B_n(t))_{n \geq 0}$ defined by the following formula

$$B_0(t) = 0, \quad B_1(t) = 1, \quad B_n(t) = \begin{cases} tB_{n/2}(t) & \text{for } n \text{ even,} \\ B_{(n+1)/2}(t) + B_{(n-1)/2}(t) & \text{for } n \text{ odd.} \end{cases}$$

This sequence was introduced in [2] as a polynomial analogue of the STERN [5] sequence $(s_n)_{n \in \mathbb{N}}$, where $s_n = B_n(1)$. Arithmetic properties of the Stern polynomials and the sequence of degrees of Stern polynomials $e(n) = \deg B_n(t)$ were considered in [6], [7]. Reducibility properties of Stern polynomials were considered in [4]. The aim of this paper is to resolve some open problems and conjectures given in [6].

In Section 2 we prove that the only possible rational roots of the Stern polynomials are in the set $\{0, -1, -\frac{1}{2}, -\frac{1}{3}\}$. We prove that each of those numbers

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is a root of infinitely many Stern polynomials. We also give some characterization of n such that $B_n(t)$ has one of these roots.

It was proven in [6] that there are no four consecutive Stern polynomials with equal degrees. In Section 3 we characterize the set of n for which $e(n) = e(n+1)$, i.e. those n for which two consecutive Stern polynomials has the same degree. We also characterize the set of n for which $e(n) = e(n+1) = e(n+2)$.

In Section 4 we investigate reciprocal Stern polynomials, i.e. polynomials such that $B_n(t) = t^{e(n)} B_n(\frac{1}{t})$. It was observed in [6] that $B_n(t)$ is reciprocal for each n of the form $2^m - 1$ or $2^m - 5$. We give two infinite sequences $(u_n), (v_n)$ such that $B_n(t)$ is reciprocal for each n of the forms $2^m - u_n$, and $2^m - v_n$.

For $n \in \mathbb{N}$ we denote by $\text{bin}(n)$ the sequence of its binary digits without leading zeros. Let $w = (w_0, w_1, \dots, w_n)$ be a finite sequence of zeros and ones, we denote by $\bar{w} = \overline{w_0 w_1 \dots w_n}$ the number $w_0 \cdot 2^n + w_1 \cdot 2^{n-1} + \dots + w_n$. For a finite sequence w of zeros and ones, we denote by w^k concatenation of k copies of w . For a set $A \subset \mathbb{Z}$ and an integer n , we denote by $A + n$ the set $\{a + n \mid a \in A\}$. For a subset $A \subset \mathbb{N}$ of non-negative integers we say that it has lower density δ if $\liminf_{n \rightarrow \infty} \frac{\#(A \cap [1, n])}{n} = \delta$. We say that it has upper density δ if $\limsup_{n \rightarrow \infty} \frac{\#(A \cap [1, n])}{n} = \delta$. If lower and upper density are both equal δ then we say that our set has density δ . The symbol $\log$ means always binary logarithm. We use standard Landau symbols.

2. Rational roots of Stern polynomials

We prove the following theorem, which can be found as a conjecture in [6, Conjecture 6.1].

Theorem 2.1. *If $a \in \mathbb{Q}$ and there exists a positive integer n such that $B_n(a) = 0$ then $a \in \{0, -1, -\frac{1}{2}, -\frac{1}{3}\}$.*

PROOF. Let $k \geq 4$ be an integer. We define a sequence $\{b_n\}_{n=0}^\infty$ by the formula $b_n = B_n(-\frac{1}{k})$. We prove that the following inequality holds

$$b_{2n+1} > \frac{1}{2} \max\{|b_n|, |b_{n+1}|\} > 0 \quad (1)$$

for every integer $n \geq 0$. In order to prove (1) we use induction. For $n = 0$ we get $1 = b_1 > \frac{1}{2} \max\{|b_0|, |b_1|\} = \frac{1}{2} > 0$. Suppose that the statement is true for all $m < n$. Let us consider two cases $n = 2s + 1$ and $n = 2s$.

If $n = 2s$ then we have

$$\begin{aligned} b_{4s+1} &= b_{2s} + b_{2s+1} = -\frac{1}{k}b_s + b_{2s+1} \geq b_{2s+1} - \frac{1}{4}|b_s| > b_{2s+1} \\ -\frac{1}{2}b_{2s+1} &= \frac{1}{2}b_{2s+1}. \end{aligned} \quad (2)$$

Moreover, $b_{2s+1} \geq \frac{1}{2}|b_s|$ and thus $\max\{|b_{2s}|, |b_{2s+1}|\} = \max\{\frac{1}{k}|b_s|, |b_{2s+1}|\} = b_{2s+1} > 0$. So in this case our inequality is proved.

If $n = 2s + 1$ then we have

$$\begin{aligned} b_{4s+3} &= b_{2s+1} + b_{2s+2} = b_{2s+1} - \frac{1}{k}b_{s+1} \geq b_{2s+1} - \frac{1}{4}|b_{s+1}| > b_{2s+1} \\ -\frac{1}{2}b_{2s+1} &= \frac{1}{2}b_{2s+1}. \end{aligned} \quad (3)$$

Moreover, $b_{2s+1} \geq \frac{1}{2}|b_{s+1}|$ and thus

$$\max\{|b_{2s+1}|, |b_{2s+2}|\} = \max\left\{\frac{1}{k}|b_{s+1}|, |b_{2s+1}|\right\} = b_{2s+1} > 0.$$

So in this case our inequality is also proved.

Now let $a \in \mathbb{Q}$ be a root of the Stern polynomial $B_n(t)$. If $a \neq 0$ then we can assume that n is odd. Moreover $B_{2s+1}(0) = 1$ for each integers $s \geq 0$, so a is of the form $\pm \frac{1}{k}$ for some integer k . But all coefficients of Stern polynomials are non-negative, so $a = -\frac{1}{k}$. Now our theorem follows from inequality (1). $\square$

Now for each $a \in \{0, -1, -\frac{1}{2}, -\frac{1}{3}\}$ we want to characterize those n for which $B_n(a) = 0$. Let us denote $R_a = \{n \in \mathbb{N} \mid B_n(a) = 0\}$. One can check that $B_n(0) = n \pmod{2}$ and $B_n(-1) = (n+1) \pmod{3} - 1$, so in this case our problem is easy and we get $R_0 = 2\mathbb{N}$ and $R_{-1} = 3\mathbb{N}$. The problem is much more complicated when $a \in \{-\frac{1}{2}, -\frac{1}{3}\}$.

Theorem 2.2. *The sets $R_{-\frac{1}{2}}, R_{-\frac{1}{3}}$ satisfies the following properties*

- (a) $2R_{-\frac{1}{2}} \subset R_{-\frac{1}{2}}, 2R_{-\frac{1}{3}} \subset R_{-\frac{1}{3}},$
- (b) $R_{-\frac{1}{2}} \subset 5\mathbb{N}, R_{-\frac{1}{3}} \subset 21\mathbb{N},$
- (c) *If $5n \in R_{-\frac{1}{2}}$, and $2^k > 5n$ then $5(2^k \pm n) \in R_{-\frac{1}{2}}$.*
- (d) *If $21n \in R_{-\frac{1}{3}}$, and $2^k > 21n$ then $21(2^k \pm n) \in R_{-\frac{1}{3}}$.*

PROOF. Part (a) is obvious. We can look only for odd elements of $R_{-\frac{1}{2}}, R_{-\frac{1}{3}}$. Looking at the reductions modulo 5 we get $B_n(-\frac{1}{2}) \equiv B_n(2) = n \pmod{5}$ so $R_{-\frac{1}{2}} \subset 5\mathbb{N}$. By looking modulo 7 we get $B_n(-\frac{1}{3}) \equiv B_n(2) = n \pmod{7}$. By

looking modulo 2 we get $B_n(-\frac{1}{3}) \equiv B_n(1) \pmod{2}$. One can easily prove that $B_n(1)$ is even if and only if n is divisible by 3, so $R_{-\frac{1}{3}} \subset 21\mathbb{N}$. Let us prove part (c). We prove by induction on n that $B_{5 \cdot 2^k \pm n}(-\frac{1}{2}) = \mp \frac{1}{4} B_n(-\frac{1}{2})$, when $2^k \geq n$. For $n = 1$ then we have to prove $B_{5 \cdot 2^k \pm 1}(-\frac{1}{2}) = \mp \frac{1}{4}$, which can be easily proven by induction on k . Now if $n = 2s$ then

$$B_{5 \cdot 2^k \pm 2s} \left(-\frac{1}{2} \right) = -\frac{1}{2} B_{5 \cdot 2^{k-1} \pm s} \left(-\frac{1}{2} \right) = -\frac{1}{2} \left(\mp \frac{1}{4} B_s \left(-\frac{1}{2} \right) \right) = \mp \frac{1}{4} B_{2s} \left(-\frac{1}{2} \right).$$

If $n = 2s + 1$ then

$$\begin{aligned} B_{5 \cdot 2^k \pm (2s+1)} \left(-\frac{1}{2} \right) &= B_{5 \cdot 2^{k-1} \pm s} \left(-\frac{1}{2} \right) + B_{5 \cdot 2^{k-1} \pm (s+1)} \left(-\frac{1}{2} \right) \\ &= \mp \frac{1}{4} \left(B_s \left(-\frac{1}{2} \right) + B_{s+1} \left(-\frac{1}{2} \right) \right) = \mp \frac{1}{4} B_{2s+1} \left(-\frac{1}{2} \right). \end{aligned}$$

Which completes the induction step, and part (c) follows. In the same manner one can prove that for $2^k > n$ we have $B_{21 \cdot 2^k \pm n}(-\frac{1}{3}) = \mp \frac{1}{27} B_n(-\frac{1}{3})$ which implies part (d). $\square$

Using previous theorem we can check that no other reductions of our sequences are eventually periodic.

Corollary 2.3. (a) For each prime number $p \neq 2, 5$ the sequence $(B_n(-\frac{1}{2}) \pmod{p})_{n \geq 0}$ is not eventually periodic.

(b) For each prime number $p \neq 2, 3, 7$ the sequence $(B_n(-\frac{1}{3}) \pmod{p})_{n \geq 0}$ is not eventually periodic.

PROOF. Let us prove part (a). Suppose that for $n > N$ we have $B_n(-\frac{1}{2}) \equiv B_{n+2^u(2v+1)}(-\frac{1}{2}) \pmod{p}$. Let us take n such that $B_{5n}(-\frac{1}{2}) \equiv 0 \pmod{p}$. We choose k such that $2^k - 1 = (2v+1)s$ for some integer s , and $2^{k+u} > 5 \cdot 2^u n$. From previous theorem we have that $B_{5(2^u n + 2^{k+u})}(-\frac{1}{2}) \equiv 0 \pmod{p}$, so $0 \equiv B_{5(2^u n + 2^{k+u}) - 5(2^u(2v+1)s)}(-\frac{1}{2}) \equiv B_{5(2^u n + 2^u)}(-\frac{1}{2}) \equiv B_{5(n+1)}(-\frac{1}{2}) \pmod{p}$, hence all numbers M divisible by 5 greater than N satisfy $B_M(-\frac{1}{2}) \equiv 0 \pmod{p}$. But one can check that $B_{15+5 \cdot 2^n}(-\frac{1}{2}) = -\frac{5}{32}$ for $n > 3$ and this number is not 0 $\pmod{p}$. We get a contradiction. Proof of part (b) can be done in the same manner. $\square$

Using reductions modulo prime numbers we prove that the sets $R_{-\frac{1}{2}}$, $R_{-\frac{1}{3}}$ have lower density 0. Before we state the next result we recall definition of a 2-automatic sequence [1],

Definition 2.4. A 2-automaton consists of

- a finite set $S = \{s_1 = I, s_2, \dots, s_n\}$, which is called the set of states. One of the states is denoted by I and called the initial state,
- two maps from S to itself F_0, F_1 ,
- a map (the output function), say π from S to a set Y .

The automaton generates a sequence $(u_n)_{n \geq 0}$ with values in Y (called a 2-automatic sequence) as follows: to compute the term u_n one expands n in base 2, say $n = \sum_{j=0}^k n_j 2^j$, then u_n is defined by $u_n = \pi(F_{n_k}(F_{n_{k-1}}(\dots(F_{n_0}(I))\dots)))$.

Now we are ready to prove the following

Theorem 2.5. Let $p > 3$ be a prime number, t be a fixed constant from the set $\{-\frac{1}{2}, -\frac{1}{3}\}$, and $d_{n,p} = \frac{\#\{0 \leq k < n \mid B_k(t) \equiv 0 \pmod{p}\}}{n}$. We have that

$$\lim_{n \rightarrow \infty} \frac{d_{2^0,p} + d_{2^1,p} + \dots + d_{2^n,p}}{n+1} \leq \frac{2}{\log p}.$$

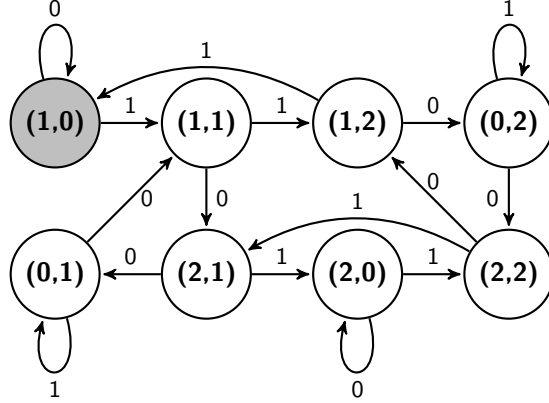
PROOF. We denote by b_n the sequence $B_n(t) \pmod{p}$. Let us consider two operators $E_0, E_1 : (\mathbb{N} \rightarrow \mathbb{F}_p) \rightarrow (\mathbb{N} \rightarrow \mathbb{F}_p)$ defined by formulas $E_0((a_n)_{n \geq 0}) = (a_{2n})_{n \geq 0}$, and $E_1((a_n)_{n \geq 0}) = (a_{2n+1})_{n \geq 0}$. Let us observe that the set $\{(\alpha b_n + \beta b_{n+1})_{n \geq 0} \mid \alpha, \beta \in \mathbb{F}_p\}$ is closed under action of E_0, E_1 because

$$E_0((\alpha b_n + \beta b_{n+1})) = (\alpha b_{2n} + \beta b_{2n+1}) = ((t\alpha + \beta)b_n + \beta b_{n+1})$$

and

$$E_1((\alpha b_n + \beta b_{n+1})) = (\alpha b_{2n+1} + \beta b_{2n+2}) = (\alpha b_n + (\alpha + t\beta)b_{n+1}).$$

Therefore (b_n) has a finite orbit under actions of E_0, E_1 so our sequence is 2-automatic. We construct automaton with states labeled with (α, β) for each $\alpha, \beta \in \mathbb{F}_p$, and edges $(\alpha, \beta) \xrightarrow{0} (t\alpha + \beta, \beta)$ and $(\alpha, \beta) \xrightarrow{1} (\alpha, t\beta + \alpha)$. The initial state is $(1, 0)$. For each state (α, β) we define its projection $\pi(\alpha, \beta) = \beta$. One can observe that if $\text{bin}(k) = w_0 w_1 \dots w_s$ then $b_k = ((E_{w_0} \circ E_{w_1} \circ \dots \circ E_{w_s})(b_n))_0$, so to compute b_k we start at initial state then go through the path $w_s w_{s-1} \dots w_0$ and finally use projection function.

Automaton for $t = -\frac{1}{2}$ and $p = 3$.

We call a state reachable if it is reachable from initial state. Let us observe that the state $(0, 1)$ is reachable. We have

$$(1, 0) \xrightarrow{1} (1, 1) \xrightarrow{0} (t+1, 1) \xrightarrow{0} (t^2+t+1, 1) \xrightarrow{0} \dots \xrightarrow{0} (t^{p-2} + \dots + t+1, 1) = (0, 1).$$

Now we can see that the state (α, β) is reachable if and only if the state (β, α) is reachable (just go to $(0, 1)$ and then use opposite arrows). Now observe that if for some $u \in \mathbb{F}_p^*$ the state $(u, 0)$ is reachable then the following different states are also reachable

$$(0, u) \xrightarrow{0} (u, u) \xrightarrow{0} (u(t+1), u) \xrightarrow{0} \dots \xrightarrow{0} (u(t^{\text{ord}(t)} - 1 + \dots + t + 1), u)$$

Where $\text{ord}(t)$ is the order of t in group $\mathbb{F}_p^*$. It is not hard to see that $\text{ord}(t) \geq \log_3 p \geq \frac{1}{2} \log p$, whence at most $\frac{2K}{\log p}$ states are of the form $(u, 0)$. Here K denotes the number of those states which are reachable from the initial state. Moreover from each reachable state there is a path to initial state. Indeed, from state (α, β) we can go to $\begin{pmatrix} t & 1 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$ or to $\begin{pmatrix} 1 & 0 \\ 1 & t \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$, those matrices has finite order in $GL_2(\mathbb{F}_p)$ so we can go backwards.

Let M be adjacency matrix of strongly connected component of our automaton as a directed graph containing initial vertex. Let $A = \frac{1}{2}M$. Using a straightforward application of Frobenius–Perron theorem (see for example [3, Example 8.3.2, p. 677]) we get that there is a limit

$$\lim_{k \rightarrow \infty} \frac{I + A + A^2 + \dots + A^k}{k+1} = G. \quad (4)$$

Of course $AG = GA = G$. As every vertex of our strongly connected component has inner and outer degree 2, one can observe that each element of the matrix G , namely $g_{i,j}$, is an arithmetic mean of g_{i,j_0} and g_{i,j_1} where $j \xrightarrow{0} j_0$ and $j \xrightarrow{1} j_1$, and on the other hand an arithmetic mean of $g_{i_0,j}$ and $g_{i_1,j}$ where $i_0 \xrightarrow{0} i$ and $i_1 \xrightarrow{1} i$. Hence using strongly connectedness we get that

$$G = \frac{1}{K} \begin{pmatrix} 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \cdots & 1 \end{pmatrix},$$

where K is the cardinality of our component.

Number $d_{2^n,p}$ is the number of paths of length n from initial state to one of states of the form $(u, 0)$ divided by 2^n . We know that A_{s_1,s_2}^n is the number of paths of length n from s_1 to s_2 , divided by 2^n . Therefore from equation (4) we get that

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{d_{0,p} + d_{2^0,p} + d_{2^1,p} + \cdots + d_{2^n,p}}{n+1} \\ = \frac{\#\{\text{reachable states of the form } (u, 0)\}}{K} \leq \frac{2}{\log p}. \end{aligned}$$

Our theorem is proved. $\square$

The above result implies the following

Corollary 2.6. *The sets $R_{-\frac{1}{2}}, R_{-\frac{1}{3}}$ has lower density 0.*

PROOF. Let t be a fixed constant from the set $\{-\frac{1}{2}, -\frac{1}{3}\}$. Let

$$d_n = \frac{\#\{0 \leq k < n \mid B_k(t) = 0\}}{n}.$$

Of course $d_{n,p} \geq d_n$ for each p . From our lemma we get

$$\limsup_{n \rightarrow \infty} \frac{d_{2^0} + d_{2^1} + \cdots + d_{2^n}}{n+1} \leq \frac{2}{\log p},$$

therefore

$$\lim_{n \rightarrow \infty} \frac{d_{2^0} + d_{2^1} + \cdots + d_{2^n}}{n+1} = 0,$$

and finally

$$\liminf_{n \rightarrow \infty} d_n = 0. \quad \square$$

We expect that $R_{-\frac{1}{2}}, R_{-\frac{1}{3}}$ has upper density 0.

Conjecture 2.7. *The sets $R_{-\frac{1}{2}}, R_{-\frac{1}{3}}$ has density 0.*

3. Consecutive polynomials with equal degrees

In this section we will characterize those n for which $e(n) = e(n+1)$. Resolving conjecture posted in [6, Conjecture 6.3]. We also characterize those n for which $e(n) = e(n+1) = e(n+2)$. We recall the recurrence satisfied by the sequence $(e_n)_{n \in \mathbb{N}}$ of degrees of Stern Polynomials

$$\begin{cases} e(0) = 1 \\ e(2n) = e(n) + 1 \\ e(4n+1) = e(n) + 1 \\ e(4n+3) = e(n+1) + 1. \end{cases}$$

Let us define sequences $(p_n)_{n \in \mathbb{N}}, (q_n)_{n \in \mathbb{N}}$ as follows

$$p_n = \frac{4^n - 1}{3}, \quad q_n = \frac{5 \cdot 4^n - 2}{3}, \quad n \geq 0.$$

Theorem 3.1. *We have the following equality*

$$\begin{aligned} \{n \mid e(n) = e(n+1)\} &= \{2^{2k+1}n + p_k \mid k, n \in \mathbb{N}_+\} \\ &\cup \{2p_k \mid k \in \mathbb{N}_+\} \cup \{2^{2k+1}n + q_k \mid k, n \in \mathbb{N}, k > 0\}. \end{aligned}$$

PROOF. We will use binary representations of integers. We write $\text{bin}(m)$ for sequence of digits of m in base two. Let us observe that $\text{bin}(p_k) = 1(01)^{k-1}$, and $\text{bin}(q_k) = 1(10)^k$ for each $k > 0$. We consider eight cases depending on the binary expansion of n .

(i) If $n = 4k$ then we have

$$e(n) = e(4k) = e(2k) + 1 = e(k) + 2$$

and

$$e(n+1) = e(4k+1) = e(k) + 1.$$

In particular $e(n) \neq e(n+1)$ in this case.

(ii) If $n = 4k+3$ then we have

$$e(n) = e(4k+3) = e(k+1) + 1$$

and

$$e(n+1) = e(4k+4) = e(k+1) + 2.$$

We thus get $e(n) \neq e(n+1)$.

(iii) If $n = \overline{\text{bin}(m)0(01)^k}$, where $m, k \in \mathbb{N}_+$ then we have

$$e(n) = e(\overline{\text{bin}(m)0(01)^k}) = e(\overline{\text{bin}(m)0}) + k = e(m) + k + 1$$

and

$$\begin{aligned} e(n+1) &= e(\overline{\text{bin}(m)0(01)^{k-1}10}) = e(\overline{\text{bin}(m)00(10)^{k-2}11}) + 1 \\ &= e(\overline{\text{bin}(m)00(10)^{k-3}11}) + 2 = \dots = e(\overline{\text{bin}(m)01}) + k = e(m) + k + 1. \end{aligned}$$

Our computation implies $e(n) = e(n+1)$.

(iv) If $n = \overline{\text{bin}(m)11(01)^k}$, where $k, m \in \mathbb{N}$, $k > 0$, then we have

$$e(n) = e(\overline{\text{bin}(m)11(01)^k}) = k + e(\overline{\text{bin}(m)11}) = k + 1 + e(\overline{\text{bin}(m+1)})$$

and

$$\begin{aligned} e(n+1) &= e(\overline{\text{bin}(m)11(01)^{k-1}10}) \\ &= e(\overline{\text{bin}(m)1(10)^{k-1}11}) + 1 = e(\overline{\text{bin}(m)1(10)^{k-2}11}) + 2 \\ &= e(\overline{\text{bin}(m)111}) + k = e(\overline{\text{bin}(m+1)0}) + k + 1 = e(\overline{\text{bin}(m+1)}) + k + 2. \end{aligned}$$

We get $e(n) \neq e(n+1)$ in this case.

(v) If $n = \overline{(01)^k}$, where $k \in \mathbb{N}_+$, then we have

$$e(n) = e(\overline{(01)^k}) = e(\overline{1}) + k - 1 = k - 1$$

and

$$\begin{aligned} e(n+1) &= e(\overline{(01)^{k-1}10}) = e(\overline{(01)^{k-1}1}) + 1 = e(\overline{(01)^{k-2}011}) + 1 \\ &= e(\overline{(01)^{k-2}1}) + 2 = e(\overline{11}) + k - 1 = k, \end{aligned}$$

and thus $e(n) \neq e(n+1)$.

(vi) If $n = \overline{\text{bin}(m)00(10)^k}$, where $m, k \in \mathbb{N}_+$, then we have

$$e(n) = e(\overline{\text{bin}(m)00(10)^k}) = k + 2 + e(\overline{\text{bin}(m)})$$

and

$$\begin{aligned} e(n+1) &= e(\overline{\text{bin}(m)0(01)^k1}) = e(\overline{\text{bin}(m)0(01)^{k-1}1}) + 1 \\ &= e(\overline{\text{bin}(m)01}) + k = e(\overline{\text{bin}(m)}) + k + 1. \end{aligned}$$

Similarly as in the previous case we get $e(n) \neq e(n+1)$.

(vii) If $n = \overline{(10)^k}$ where $k \in \mathbb{N}_+$, then we have

$$e(n) = e(\overline{(10)^k}) = e(\overline{(01)^k}) + 1 = e(1) + k = k$$

and

$$e(n+1) = e(\overline{(10)^{k-1}11}) = e(\overline{(10)^{k-2}11}) + 1 = e(\overline{11}) + k - 1 = k,$$

and thus $e(n) = e(n+1)$.

(viii) If $n = \overline{\text{bin}(m)1(10)^k}$ where $m, k \in \mathbb{N}$, $k > 0$, then we have

$$e(n) = e(\overline{\text{bin}(m)1(10)^k}) = e(\overline{\text{bin}(m)11(01)^{k-1}}) + 1 = e(\overline{\text{bin}(m+1)}) + k + 1$$

and

$$\begin{aligned} e(n+1) &= e(\overline{\text{bin}(m)1(10)^{k-1}11}) = e(\overline{\text{bin}(m)1(10)^{k-2}11}) + 1 \\ &= e(\overline{\text{bin}(m)111}) + k - 1 = e(\overline{\text{bin}(m+1)}) + k + 1. \end{aligned}$$

and thus $e(n) = e(n+1)$.

All numbers of the form $4k+1$ can be written uniquely in one of the forms (iii), (iv), (v), and all numbers of the form $4k+2$ can be written uniquely in one of the forms (vi), (vii), (viii). Therefore all cases were considered. In cases (iii), (vii), (viii) we have that $e(n) = e(n+1)$, and one can observe that the numbers in this cases are exactly the elements of the set $\{2^{2k+1}n + p_k \mid k, n \in \mathbb{N}_+\} \cup \{2p_k \mid k \in \mathbb{N}_+\} \cup \{2^{2k+1}n + q_k \mid k, n \in \mathbb{N}, k > 0\}$. So our theorem is proved. $\square$

In the next theorem we characterize the set of n for which $e(n) = e(n+1) = e(n+2)$.

Theorem 3.2. *We have the following equality*

$$\begin{aligned} \{n \mid e(n) = e(n+1) = e(n+2)\} &= \{2^{2k+1}n + p_k \mid k, n \in \mathbb{N}_+, k \geq 2\} \\ &\cup \{2p_k - 1 \mid k \geq 2\} \cup \{2^{2k+1}n + q_k - 1 \mid k, n \in \mathbb{N}, k \geq 2\}. \end{aligned}$$

PROOF. Let $A = \{2^{2k+1}n + p_k \mid k, n \in \mathbb{N}_+\}$, $B = \{2p_k \mid k \in \mathbb{N}_+\}$, and $C = \{2^{2k+1}n + q_k \mid k, n \in \mathbb{N}, k > 0\}$. We want to compute

$$((A \cup B \cup C) \cap ((A \cup B \cup C) + 1)) - 1.$$

But as every element of A equals 1 (mod 4) and every element of $B \cup C$ equals 2 (mod 4), it is enough to compute $((A+1) \cap B) \cup ((A+1) \cap C)$. First let us compute $(A+1) \cap B$, we have

$$2^{2k+1}n + p_k + 1 = 2p_l, \quad \text{for } n, k, l \geq 1.$$

By a simple computations we get

$$3 \cdot 2^{2k+1}n + 4^k = 2 \cdot 4^l - 4.$$

Looking modulo 4 we get that $k = 1$ or $l = 1$. When $k = 1$ then our equality became $12n + 4 = 4^l$. For each l we have exactly one n such that this equality holds, but as $n \geq 2$ we have $l \geq 2$. If $l = 1$ then our equality became $2^{2k+1}n + p_k + 1 = 2$ which contradicts $n, k \geq 1$. Summarizing $(A + 1) \cap B = \{2p_l - 1 \mid l \geq 2\}$. Let us compute $(A + 1) \cap C$. We have

$$2^{2k+1}n + p_k + 1 = 2^{2l+1}m + q_l, \quad \text{for } n, k, l \geq 1, m \geq 0.$$

By a simple computations we get

$$2^{2k+1} \cdot 3n + 4^k = 2^{2l+1} \cdot 3m + 5 \cdot 4^l - 4.$$

Looking (mod 8) we get that $k = 1$ or $l = 1$. In the first case we get $8n + 2 = 2^{2l+1}m + q_l$. For each l, m such that $l \geq 2$ we have exactly one n such that this equality holds. So we get $\{2^{2l+1}m + q_l \mid m, l \in \mathbb{N}, l \geq 2\}$. In the second case we get $\{2^{2k+1}n + p_k + 1 \mid k, n \in \mathbb{N}_+, k \geq 2\}$ in the same way. Summarizing we get $(A+1) \cap C = \{2^{2l+1}m + q_l \mid m, l \in \mathbb{N}, l \geq 2\} \cup \{2^{2k+1}n + p_k + 1 \mid k, n \in \mathbb{N}_+, k \geq 2\}$. And our result follows. $\square$

4. Reciprocal Stern polynomials

Let us recall that the polynomial $B_n(t)$ is called reciprocal if it satisfies equality $B_n(t) = t^{e(n)}B_n(\frac{1}{t})$. It was observed in [6] that for each n polynomials $B_{2^n-1}(t)$ and $B_{2^n-5}(t)$ are reciprocal. We define sequences $(u_n)_{n \in \mathbb{N}}$ and $(v_n)_{n \in \mathbb{N}}$ as follows

$$\begin{cases} u_0 = 1 & u_n = 2^{8n-2}u_{n-1} + 2^{4n} - 1, \\ v_0 = 5 & v_n = 2^{8n+2}v_{n-1} - 2^{4n+2} + 1, \end{cases}$$

and show that whenever $2^k > u_n$ (or $2^k > v_n$) then $B_{2^k-u_n}$ (or $B_{2^k-v_n}$) is reciprocal.

Theorem 4.1. *For each integer $n \geq 0$ all polynomials in the sequences $(B_{2^k-v_n})_{k \geq 4n^2+6n+3}$, and $(B_{2^k-u_n})_{k \geq 4n^2+2n+1}$ are reciprocal.*

PROOF. Let us observe that $\text{bin}(u_n) = 10^2 1^4 0^6 \dots 0^{4n-2} 1^{4n}$ and $\text{bin}(v_n) = 10^2 1^4 \dots 1^{4n} 0^{4n+1} 1$. Using recursive formula for $e(n)$, binary representations and induction we will get that $e(2^k - u_n) = k - 2n - 1$ and $e(2^k - v_n) = k - 2n - 2$. Let

us prove our theorem. We proceed by induction. For $u_0 = 1$ we have $B_{2^k-1} = t^{k-1} + t^{k-2} + \dots + t + 1$ and our conditions are satisfied. For $v_0 = 5$ we have

$$\begin{aligned} B_{2^k-5}(t) &= (1+t)B_{2^k-2-1}(t) + tB_{2^k-3-1}(t) \\ &= (1+t)t^{k-3}B_{2^k-2-1}(t^{-1}) + t^{k-3}B_{2^k-3-1}(t^{-1}) \\ &= t^{k-2}((1+t^{-1})B_{2^k-2-1}(t^{-1}) + t^{-1}B_{2^k-3-1}(t^{-1})) = t^{k-2}B_{2^k-5}(t^{-1}). \end{aligned}$$

Let us compute

$$\begin{aligned} B_{2^k-u_n}(t) &= B_{2^k-\overline{10^2 1^4 0^6 \dots 0^{4n-2} 1^{4n}}}(t) \\ &= B_{2^k-1-\overline{10^2 1^4 0^6 \dots 0^{4n-2} 1^{4n-1}}}(t) + B_{2^k-1-\overline{10^2 1^4 0^6 \dots 0^{4n-3} 10^{4n-1}}}(t) \\ &= B_{2^k-1-\overline{10^2 1^4 0^6 \dots 0^{4n-2} 1^{4n-1}}}(t) + t^{4n-1}B_{2^k-4n-\overline{10^2 1^4 0^6 \dots 0^{4n-3} 1}}(t) \\ &= B_{2^k-1-\overline{10^2 1^4 0^6 \dots 0^{4n-2} 1^{4n-1}}}(t) + t^{4n-1}B_{2^k-4n-v_{n-1}}(t) \\ &= \dots \\ &= B_{2^k-4n-\overline{10^2 1^4 0^6 \dots 0^{4n-2}}}(t) + (t^{4n-1} + \dots + t + 1)B_{2^k-4n-v_{n-1}}(t) \\ &= t^{4n-2}B_{2^k-8n+2-u_{n-1}}(t) + (t^{4n-1} + \dots + t + 1)B_{2^k-4n-v_{n-1}}(t) \\ &= (*) \end{aligned}$$

Using induction hypothesis we get

$$\begin{aligned} (*) &= t^{k-6n+1}B_{2^k-8n+2-u_{n-1}}(t^{-1}) + (t^{4n-1} + \dots + t + 1)t^{k-6n}B_{2^k-4n-v_{n-1}}(t^{-1}) \\ &= t^{k-2n-1} \left(t^{2-4n}B_{2^k-8n+2-u_{n-1}}(t^{-1}) + \left(\frac{t^{4n}-1}{t-1} \right) t^{1-4n}B_{2^k-4n-v_{n-1}}(t^{-1}) \right) \\ &= t^{k-2n-1}B_{2^k-u_n}(t^{-1}). \end{aligned}$$

So $B_{2^k-u_n}$ is reciprocal. In the same way one can prove that $B_{2^k-v_n}$ is reciprocal. We omit the details. $\square$

Corollary 4.2. *Let $\text{Rec} = \{n \mid B_n(t) = t^{e(n)}B_n(t^{-1})\}$. We have that $\#(\text{Rec} \cap [1, n]) = \Omega(\log(n)^{3/2})$.*

PROOF. Let $n > 2^k$ we have that

$$\begin{aligned} \#(\text{Rec} \cap [1, n]) &\geq \#\{(i, m) \mid u_m < 2^i, i < k\} = \sum_{u_m < 2^k} (k-1 - \lfloor \log(u_m) \rfloor) \\ &\sum_{4m^2+2m+1 < k} (k - (4m^2 + 2m)) = \frac{1}{2}k\sqrt{k} - \frac{4}{3} \left(\frac{\sqrt{k}}{2} \right)^3 + o(k\sqrt{k}) = \theta(k\sqrt{k}), \end{aligned}$$

and our result follows. $\square$

We expect the following conjecture about density of the set Rec

Conjecture 4.3. *The function $f(n) = \#(\text{Rec} \cap [1, n])$ is $O(\log(n)^k)$ for some constant k .*

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